

Solving Boolean Equation

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Abstract

If $\forall P$: proposition, $B(P)$ is the truth value(0 or 1) of P then we can solve a boolean equation by using these below.

$$B(p_1 \vee p_2 \vee \dots \vee p_n) \equiv 1 + \prod_{k=1}^n (1 + p_k) \pmod{2}$$

$$\begin{aligned} \left\{ (x_1, x_2, \dots, x_n) \mid \prod_{i=1}^n B(x_i) \equiv 0 \pmod{2} \right\} &= \left(\bigcap_{i=1}^n \{ (x_i) \mid B(x_i) \equiv 1 \pmod{2} \} \right)^c \\ &= \{ (x_1, x_2, \dots, x_n) \mid (1, 1, 1, \dots, 1) \}^c \end{aligned}$$

Introduction

We have started to study because we don't know how to apply many Boolean laws to Boolean algebra. So, We study to redefine Boolean algebra using by mod 2, to reprove Boolean laws, and we study the method of solving Boolean equation.

1. Redefine Boolean algebra and basical theorem

Definition 1.1) $\forall P$: proposition

- $B(P) = \begin{cases} 0, & \text{truth value of } P \text{ is } False \\ 1, & \text{truth value of } P \text{ is } True \end{cases}$
- If $B(P) = z(0 \text{ or } 1)$ then $B(P) \equiv z(mod 2)$
(But, we can omit $B()$ like as $B(P) \equiv P(mod 2)$ for usability.)

Ex) $P \equiv p \vee q(p, q : \text{proposition})$, If the truth value of P is *True*
then $B(P) \equiv B(p \vee q) \equiv P \equiv 1(mod 2)$

Theorem 1.1) $\forall n \in \mathbb{N}, \forall P$: proposition

$$\begin{cases} 2nB(P) \equiv 2nP \equiv 0(mod 2) \\ (2n + 1)B(P) \equiv B(P) \equiv (2n + 1)P \equiv P(mod 2) \end{cases}$$

$$(B(P))^n \equiv P^n \equiv B(P) \equiv P(mod 2)$$

$$-B(P) \equiv -P \equiv B(P) \equiv P(mod 2)$$

Proof)

$$\begin{cases} 2nB(P) \equiv 2 \times nB(P) \equiv 0 \times nB(P) \equiv 0(mod 2) \\ (2n + 1)B(P) \equiv 2nB(P) + B(P) \equiv 0 + B(P) \equiv B(P)(mod 2) \end{cases}$$

$$\begin{cases} B(P) \equiv 1(mod 2) \Rightarrow (B(P))^n \equiv (1)^n \equiv 1 \equiv B(P)(mod 2) \\ B(P) \equiv 0(mod 2) \Rightarrow (B(P))^n \equiv (0)^n \equiv 0 \equiv B(P)(mod 2) \end{cases}$$

$$\begin{cases} B(P) \equiv 1(mod 2) \Rightarrow -B(P) \equiv -1 \times 1 \equiv -1 \equiv 0 - 1 \equiv 2 - 1 \equiv 1 \equiv B(P)(mod 2) \\ B(P) \equiv 0(mod 2) \Rightarrow -B(P) \equiv -1 \times 0 \equiv 0 \equiv B(P)(mod 2) \end{cases} \blacksquare$$

Theorem 1.2) $\forall p, q$: proposition

$$B(c) \equiv 0(mod 2), B(t) \equiv 1(mod 2)(\text{where, } c : \text{contradiction, } t : \text{tautology})$$

$$B(\sim p) \equiv 1 + B(p) \equiv 1 + p(mod 2)$$

$$B(p \wedge q) \equiv B(p)B(q) \equiv pq(mod 2)$$

$$B(p \vee q) \equiv B(p) + B(q) + B(p)B(q) \equiv p + q + pq(mod 2)$$

$$B(p \rightarrow q) \equiv 1 + B(p) + B(p)B(q) \equiv 1 + p + pq(mod 2)$$

$$B(p \leftrightarrow q) \equiv 1 + B(p) + B(q) \equiv 1 + p + q(mod 2)$$

Proof)

p	q	$\sim p$	$p \wedge q$	$p \vee q$	$p \rightarrow q$	$p \leftrightarrow q$
		$1 + p$	pq	$p + q + pq$	$1 + p + pq$	$1 + p + q$
0	0	1	0	0	1	1
0	1	1	0	1	1	0
1	0	0	0	1	0	0
1	1	0	1	1	1	1

■

Theorem 1.3) $\forall p, q$: proposition

$$B(p \text{ XOR } q) \equiv B(p) + B(q) \equiv p + q \pmod{2}$$

$$B(p \text{ NAND } q) \equiv 1 + B(p)B(q) \equiv 1 + pq \pmod{2}$$

$$B(p \text{ NOR } q) \equiv 1 + B(p) + B(q) + B(p)B(q) \equiv 1 + p + q + pq \pmod{2}$$

$$B(p \text{ XNOR } q) \equiv 1 + B(p) + B(q) \equiv 1 + p + q \pmod{2}$$

Proof)

p	q	$p \text{ XOR } q$	$p \text{ NAND } q$	$p \text{ NOR } q$	$p \text{ XNOR } q$
		$p + q$	$1 + pq$	$1 + p + q + pq$	$1 + p + q$
0	0	0	1	1	1
0	1	1	1	0	0
1	0	1	1	0	0
1	1	0	0	0	1

■

Corollary 1.1) $\forall p, q$: proposition

$$B(p \Rightarrow q) \Leftrightarrow p + pq \equiv 0 \pmod{2}$$

$$B(p \Leftrightarrow q) \Leftrightarrow p + q \equiv 0 \pmod{2}$$

Proof)

$$B(p \Rightarrow q) \Leftrightarrow B(p \rightarrow q) \equiv 1 \pmod{2} \Leftrightarrow 1 + p + pq \equiv 1 \pmod{2} \Leftrightarrow p + pq \equiv 0 \pmod{2}$$

$$B(p \Leftrightarrow q) \Leftrightarrow B(p \leftrightarrow q) \equiv 1 \pmod{2} \Leftrightarrow 1 + p + q \equiv 1 \pmod{2} \Leftrightarrow p + q \equiv 0 \pmod{2} \blacksquare$$

Corollary 1.2) $\forall p, q$: proposition,

$$B(p \vee q) \equiv p + q + pq \equiv 1 + (1 + p)(1 + q) \pmod{2}$$

$$B(p \rightarrow q) \equiv 1 + p + pq \equiv 1 + p(1 + q) \pmod{2}$$

$$B(p \leftrightarrow q) \equiv 1 + p + q \equiv 1 + (1 + pq)(p + q + pq) \equiv (1 + p + pq)(1 + q + qp) \pmod{2}$$

$$B(p \text{ XOR } q) \equiv p + q \equiv (1 + pq)(p + q + pq) \equiv 1 + (1 + p + pq)(1 + q + qp) \pmod{2}$$

Proof)

$$\begin{aligned} B(p \vee q) &\equiv p + q + pq \equiv 1 + 1 + p + q + pq \equiv 1 + p(1 + q) + 1 + q \\ &\equiv 1 + (1 + p)(1 + q) \pmod{2} \end{aligned}$$

$$B(p \rightarrow q) \equiv 1 + p + pq \equiv 1 + p(1 + q) \pmod{2}$$

$$B(p \leftrightarrow q) \equiv 1 + p + q \equiv 1 + p + q + pq + pq \equiv (1 + pq) + (p + q + pq)$$

$$\equiv (1 + pq) + (p + q + pq) + ppq + qpq + pqpq + pq$$

$$\equiv (1 + pq) + (p + q + pq) + (p + q + pq)pq + pq$$

$$\equiv (1 + pq) + (p + q + pq)(1 + pq) + pq$$

$$\equiv (1 + pq + pq) + (p + q + pq)(1 + pq) \equiv 1 + (1 + pq)(p + q + pq)$$

$$\equiv (1 + p + pq) + q(1 + p + pq) + pq(1 + p + pq)$$

$$\equiv (1 + p + pq)(1 + q + pq) \pmod{2}$$

$$B(p \text{ XOR } q) \equiv p + q \equiv 1 + 1 + p + q \equiv 1 + B(p \leftrightarrow q) \equiv (1 + pq)(p + q + pq)$$

$$\equiv 1 + (1 + p + pq)(1 + q + pq) \pmod{2} \blacksquare$$

Corollary 1.3) $\forall p_k$: proposition,

$$1) B(p_1 \wedge p_2 \wedge \dots \wedge p_n) \equiv \prod_{k=1}^n p_k \pmod{2}$$

$$2) B(p_1 \vee p_2 \vee \dots \vee p_n) \equiv 1 + \prod_{k=1}^n (1 + p_k) \pmod{2}$$

Proof)

$$1) B(p_1 \wedge p_2) \equiv \prod_{k=1}^2 p_k \pmod{2} \text{ by Theorem 1.2}$$

$$\text{Assume } B(p_1 \wedge p_2 \wedge \dots \wedge p_n) \equiv \prod_{k=1}^n p_k \pmod{2} \Rightarrow$$

$$\begin{aligned} B(p_1 \wedge p_2 \wedge \dots \wedge p_n \wedge p_{n+1}) &\equiv B((p_1 \wedge p_2 \wedge \dots \wedge p_n) \wedge p_{n+1}) \equiv \left(\prod_{k=1}^n p_k \right) p_{n+1} \\ &\equiv \prod_{k=1}^{n+1} p_k \pmod{2} \end{aligned}$$

$$2) B(p_1 \vee p_2) \equiv 1 + \prod_{k=1}^2 (1 + p_k) \pmod{2} \text{ by Corollary 1.2}$$

$$\text{Assume } B(p_1 \vee p_2 \vee \dots \vee p_n) \equiv 1 + \prod_{k=1}^n (1 + p_k) \pmod{2} \Rightarrow$$

$$\begin{aligned} B(p_1 \vee p_2 \vee \dots \vee p_n \vee p_{n+1}) &\equiv B((p_1 \vee p_2 \vee \dots \vee p_n) \vee p_{n+1}) \\ &\equiv 1 + \left(1 + \left(1 + \prod_{k=1}^n (1 + p_k) \right) \right) (1 + p_{n+1}) \\ &\equiv 1 + \left(1 + 1 + \prod_{k=1}^n (1 + p_k) \right) (1 + p_{n+1}) \equiv 1 + \left(\prod_{k=1}^n (1 + p_k) \right) (1 + p_{n+1}) \\ &\equiv 1 + \left(\prod_{k=1}^{n+1} (1 + p_k) \right) \pmod{2} \blacksquare \end{aligned}$$

2. Redefine Boolean laws

2.1) $p \rightarrow q \equiv \sim(p \wedge \sim q)$

Proof)

$$B(p \rightarrow q) \equiv 1 + p + pq \pmod{2}$$

$$B(\sim(p \wedge \sim q)) \equiv 1 + B(p \wedge \sim q) \equiv 1 + p(1 + q) \equiv 1 + p + pq \pmod{2} \blacksquare$$

2.2) $p \leftrightarrow q \equiv (p \rightarrow q) \wedge (q \rightarrow p)$

Proof)

$$B(p \leftrightarrow q) \equiv 1 + p + q \pmod{2}$$

$$B((p \rightarrow q) \wedge (q \rightarrow p)) \equiv B(p \rightarrow q)B(q \rightarrow p) \equiv (1 + p + pq)(1 + q + qp)$$

$$\equiv 1 + q + qp + p + pq + p^2q + pq + pq^2 + p^2q^2$$

$$\equiv 1 + q + qp + p + pq + pq + pq + pq \equiv 1 + q + p + 6pq$$

$$\equiv 1 + p + q \pmod{2} \blacksquare$$

2.3) $\sim(\sim p) \equiv p$ (Law of Double Negation)

Proof)

$$B(\sim(\sim p)) \equiv 1 + B(\sim p) \equiv 1 + (1 + p) \equiv 2 + p \equiv p \pmod{2} \blacksquare$$

2.4) $p \wedge p \equiv p$, $p \vee p \equiv p$ (Laws of Idempotency)

Proof)

$$1) B(p \wedge p) \equiv pp \equiv p^2 \equiv p \pmod{2} \blacksquare$$

$$2) B(p \vee p) \equiv p + p + pp \equiv p + p + p^2 \equiv p + p + p \equiv 3p \equiv p \pmod{2} \blacksquare$$

2.5) $p \rightarrow q \equiv \sim q \rightarrow \sim p$ (Contrapositive Law)

Proof)

$$B(p \rightarrow q) \equiv 1 + p + pq \pmod{2}$$

$$B(\sim q \rightarrow \sim p) \equiv 1 + B(\sim q) + B(\sim q)B(\sim p) \equiv 1 + (1 + q) + (1 + q)(1 + p)$$

$$\equiv 1 + (1 + q) + (1 + p + q + qp) \equiv 3 + 2q + p + pq$$

$$\equiv 1 + p + pq \pmod{2} \blacksquare$$

2.6) $\sim(p \wedge q) \equiv \sim p \vee \sim q$, $\sim(p \vee q) \equiv \sim p \wedge \sim q$ (De Morgan's Laws)

Proof)

$$1) B(\sim(p \wedge q)) \equiv 1 + B(p \wedge q) \equiv 1 + pq \pmod{2}$$

$$B(\sim p \vee \sim q) \equiv B(\sim p) + B(\sim q) + B(\sim p)B(\sim q) \equiv (1 + p) + (1 + q) + (1 + p)(1 + q)$$

$$\equiv (1 + p) + (1 + q) + (1 + p + q + qp) \equiv 3 + 2p + 2q + qp$$

$$\equiv 1 + pq \pmod{2} \blacksquare$$

$$2) B(\sim(p \vee q)) \equiv 1 + B(p \vee q) \equiv 1 + (p + q + pq) \pmod{2}$$

$$B(\sim p \wedge \sim q) \equiv B(\sim p)B(\sim q) \equiv (1 + p)(1 + q) \equiv 1 + q + p + pq \pmod{2} \blacksquare$$

2.7) $p \wedge q \equiv q \wedge p, p \vee q \equiv q \vee p$ (Commutative Laws)

Proof)

$$1) B(p \wedge q) \equiv pq(\text{mod } 2), B(q \wedge p) \equiv qp(\text{mod } 2) \blacksquare$$

$$2) B(p \vee q) \equiv p + q + pq(\text{mod } 2), B(q \vee p) \equiv q + p + qp(\text{mod } 2) \blacksquare$$

2.8) $(p \wedge q) \wedge r \equiv p \wedge (q \wedge r), (p \vee q) \vee r \equiv p \vee (q \vee r)$ (Associative Laws)

Proof)

$$1) B((p \wedge q) \wedge r) \equiv (pq)r \equiv pqr(\text{mod } 2), B(p \wedge (q \wedge r)) \equiv p(qr) \equiv pqr(\text{mod } 2) \blacksquare$$

$$\begin{aligned} 2) B((p \vee q) \vee r) &\equiv B(p \vee q) + r + B(p \vee q)r \equiv B(p \vee q)(1 + r) + r \\ &\equiv (p + q + pq)(1 + r) + r \equiv (p + q + pq)(1 + r) + r + 1 + 1 \\ &\equiv (p + q + pq + 1)(1 + r) + 1 \equiv ((p + 1) + q(1 + p))(1 + r) + 1 \\ &\equiv (p + 1)(1 + q)(1 + r) + 1(\text{mod } 2) \end{aligned}$$

$$\begin{aligned} B(p \vee (q \vee r)) &\equiv p + B(q \vee r) + pB(q \vee r) \equiv p + (1 + p)B(q \vee r) \\ &\equiv p + (1 + p)(q + r + qr) \equiv 1 + 1 + p + (1 + p)(q + r + qr) \\ &\equiv 1 + (1 + p)(1 + q + r + qr) \equiv 1 + (1 + p)(1 + r + q(1 + r)) \\ &\equiv 1 + (1 + p)(1 + q)(1 + r)(\text{mod } 2) \blacksquare \end{aligned}$$

2.9) $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r),$

$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$ (Distributive Laws)

Proof)

$$1) B(p \wedge (q \vee r)) \equiv pB(q \vee r) \equiv p(q + r + qr) \equiv pq + pr + pqr(\text{mod } 2)$$

$$\begin{aligned} B((p \wedge q) \vee (p \wedge r)) &\equiv B(p \wedge q) + B(p \wedge r) + B(p \wedge q)B(p \wedge r) \\ &\equiv (pq) + (pr) + (pq)(pr) \equiv (pq) + (pr) + (p^2qr) \\ &\equiv (pq) + (pr) + (pqr) \equiv pq + pr + pqr(\text{mod } 2) \blacksquare \end{aligned}$$

$$2) B(p \vee (q \wedge r)) \equiv p + B(q \wedge r) + pB(q \wedge r) \equiv p + qr + pqr(\text{mod } 2)$$

$$\begin{aligned} B((p \vee q) \wedge (p \vee r)) &\equiv B(p \vee q)B(p \vee r) \equiv (p + q + pq)(p + r + pr) \\ &\equiv p^2 + qp + p^2q + pr + qr + pqr + p^2r + qpr + p^2qr \\ &\equiv p + qp + pq + pr + qr + pqr + pr + qpr + pqr \\ &\equiv p + 2pq + 2pr + qr + 3pqr \equiv p + qr + pqr(\text{mod } 2) \blacksquare \end{aligned}$$

2.10) $p \wedge t \leftrightarrow p, p \vee t \leftrightarrow t, p \vee c \leftrightarrow p, p \wedge c \leftrightarrow c, c \Rightarrow p, p \Rightarrow t$

Proof)

- 1) $B(p \wedge t \leftrightarrow p) \equiv 1 + B(p \wedge t) + B(p) \equiv 1 + (p \times 1) + p \equiv 1 + p + p \equiv 1 + 2p \equiv 1 \pmod{2}$ ■
- 2) $B(p \vee t \leftrightarrow t) \equiv 1 + B(p \vee t) + B(t) \equiv 1 + (p + 1 + p \times 1) + 1 \equiv 1 + (p + 1 + p) + 1 \equiv 3 + 2p \equiv 1 \pmod{2}$ ■
- 3) $B(p \vee c \leftrightarrow p) \equiv 1 + B(p \vee c) + B(p) \equiv 1 + (p + 0 + p \times 0) + p \equiv 1 + (p) + p \equiv 1 + 2p \equiv 1 \pmod{2}$ ■
- 4) $B(p \wedge c \leftrightarrow c) \equiv 1 + B(p \wedge c) + B(c) \equiv 1 + (p \times 0) + 0 \equiv 1 + (0) + 0 \equiv 1 \pmod{2}$ ■
- 5) $B(c \rightarrow p) \equiv 1 + B(c) + B(c)p \equiv 1 + (0) + (0 \times p) \equiv 1 + (0) + (0) \equiv 1 \pmod{2}$ ■
- 6) $B(p \rightarrow t) \equiv 1 + p + pB(t) \equiv 1 + p + (p \times 1) \equiv 1 + p + (p) \equiv 1 + 2p \equiv 1 \pmod{2}$ ■

2.11) $p \wedge \sim p \equiv c, p \vee \sim p \equiv t$

Proof)

- 1) $B(p \wedge \sim p) \equiv p(1 + p) \equiv p + p^2 \equiv p + p \equiv 2p \equiv 0 \equiv B(c) \pmod{2}$ ■
- 2) $B(p \vee \sim p) \equiv p + (1 + p) + p(1 + p) \equiv p + 1 + p + p + p^2 \equiv 1 + 3p + p \equiv 1 + 4p \equiv 1 \equiv B(t) \pmod{2}$ ■

2.12) $p \Rightarrow p \vee q$ (Law of addition)

Proof)

$$\begin{aligned} B(p \rightarrow p \vee q) &\equiv 1 + p + pB(p \vee q) \equiv 1 + p + p(p + q + pq) \equiv 1 + p + (p^2 + pq + p^2q) \\ &\equiv 1 + p + (p + pq + pq) \equiv 1 + 2p + 2pq \equiv 1 \pmod{2} \quad \blacksquare \end{aligned}$$

2.13) $p \wedge q \Rightarrow p, p \wedge q \Rightarrow q$ (Laws of Simplification)

Proof)

- 1) $B(p \wedge q \rightarrow p) \equiv 1 + B(p \wedge q) + B(p \wedge q)p \equiv 1 + pq + pq(p) \equiv 1 + pq + p^2q \equiv 1 + pq + pq \equiv 1 + 2pq \equiv 1 \pmod{2}$ ■
- 2) $B(p \wedge q \rightarrow q) \equiv 1 + B(p \wedge q) + B(p \wedge q)q \equiv 1 + pq + pq(q) \equiv 1 + pq + pq^2 \equiv 1 + pq + pq \equiv 1 + 2pq \equiv 1 \pmod{2}$ ■

2.14) $p \wedge q \rightarrow r \equiv p \rightarrow (q \rightarrow r)$ (Exportation Law)

Proof)

$$\begin{aligned} B(p \wedge q \rightarrow r) &\equiv 1 + B(p \wedge q) + B(p \wedge q)r \equiv 1 + pq + pqr \pmod{2} \\ B(p \rightarrow (q \rightarrow r)) &\equiv 1 + p + pB(q \rightarrow r) \equiv 1 + p + p(1 + q + qr) \equiv 1 + p + p + pq + pqr \\ &\equiv 1 + 2p + pq + pqr \equiv 1 + pq + pqr \pmod{2} \quad \blacksquare \end{aligned}$$

2.15) $(p \rightarrow q) \wedge (q \rightarrow r) \Rightarrow (p \rightarrow r)$ (Transitive Law)

Proof)

$$\begin{aligned}
 B((p \rightarrow q) \wedge (q \rightarrow r) \rightarrow (p \rightarrow r)) &\equiv 1 + B((p \rightarrow q) \wedge (q \rightarrow r))(1 + B(p \rightarrow r)) \\
 &\equiv 1 + (1 + p + pq)(1 + q + qr)(1 + (1 + p + pr)) \\
 &\equiv 1 + (1 + p + pq)(1 + q + qr)(2 + p + pr) \\
 &\equiv 1 + (1 + p + pq)(1 + q + qr)(p + pr) \\
 &\equiv 1 + (1 + p + pq)(1 + q + qr)(1 + r)p \\
 &\equiv 1 + (1 + p + pq + q + pq + pq^2 + qr + pqr + pq^2r)(1 + r)p \\
 &\equiv 1 + (1 + p + pq + q + pq + pq + qr + pqr + pqr)(1 + r)p \\
 &\equiv 1 + (1 + p + 3pq + q + qr + 2pqr)(1 + r)p \\
 &\equiv 1 + (1 + p + pq + q + qr)p(1 + r) \\
 &\equiv 1 + (p + p^2 + p^2q + pq + pqr)(1 + r) \\
 &\equiv 1 + (p + p + pq + pq + pqr)(1 + r) \equiv 1 + (2p + 2pq + pqr)(1 + r) \\
 &\equiv 1 + (pqr)(1 + r) \equiv 1 + (pqr + pqr^2) \equiv 1 + (pqr + pqr) \equiv 1 + (2pqr) \\
 &\equiv 1 \pmod{2} \blacksquare
 \end{aligned}$$

2.16) $(p \rightarrow q) \Leftrightarrow (p \wedge \sim q \rightarrow q \wedge \sim q)$ (Reductio ad Absurdum)

Proof)

$$\begin{aligned}
 B((p \rightarrow q) \leftrightarrow (p \wedge \sim q \rightarrow q \wedge \sim q)) &\equiv 1 + B(p \rightarrow q) + B(p \wedge \sim q \rightarrow q \wedge \sim q) \\
 &\equiv 1 + (1 + p + pq) + (1 + B(p \wedge \sim q)(1 + B(q \wedge \sim q))) \\
 &\equiv 1 + (1 + p + pq) + (1 + (p(1 + q))(1 + q(1 + q))) \\
 &\equiv 1 + (1 + p + pq) + (1 + (p + pq)(1 + q + q^2)) \\
 &\equiv 1 + (1 + p + pq) + (1 + (p + pq)(1 + q + q)) \\
 &\equiv 1 + (1 + p + pq) + (1 + (p + pq)(1 + 2q)) \\
 &\equiv 1 + (1 + p + pq) + (1 + (p + pq)) \equiv 1 + 2(1 + p + pq) \equiv 1 \pmod{2} \blacksquare
 \end{aligned}$$

2.17) $(p \vee q) \wedge \sim p \Rightarrow q$ (Disjunctive Syllogism)

Proof)

$$\begin{aligned} B((p \vee q) \wedge \sim p \rightarrow q) &\equiv 1 + B((p \vee q) \wedge \sim p)(1 + q) \equiv 1 + B(p \vee q)(1 + p)(1 + q) \\ &\equiv 1 + (p + q + pq)(1 + p)(1 + q) \\ &\equiv 1 + (p + p^2 + q + qp + pq + p^2q)(1 + q) \\ &\equiv 1 + (p + p + q + qp + pq + pq)(1 + q) \equiv 1 + (2p + q + 3pq)(1 + q) \\ &\equiv 1 + (q + pq)(1 + q) \equiv 1 + (q + q^2 + pq + pq^2) \\ &\equiv 1 + (q + q + pq + pq) \equiv 1 + (2q + 2pq) \equiv 1(\text{mod } 2) \blacksquare \end{aligned}$$

2.18) $(p \rightarrow q) \wedge p \Rightarrow q$ (Modus Ponens)

Proof)

$$\begin{aligned} B((p \rightarrow q) \wedge p \rightarrow q) &\equiv 1 + B((p \rightarrow q) \wedge p)(1 + q) \equiv 1 + ((1 + p + pq)p)(1 + q) \\ &\equiv 1 + (p + p^2 + p^2q)(1 + q) \equiv 1 + (p + p + pq)(1 + q) \\ &\equiv 1 + (2p + pq)(1 + q) \equiv 1 + pq(1 + q) \equiv 1 + pq + pq^2 \equiv 1 + pq + pq \\ &\equiv 1 + 2pq \equiv 1(\text{mod } 2) \blacksquare \end{aligned}$$

2.19) $(p \rightarrow q) \wedge \sim q \Rightarrow \sim p$ (Modus Tollens)

Proof)

$$\begin{aligned} B((p \rightarrow q) \wedge \sim q \rightarrow \sim p) &\equiv 1 + B((p \rightarrow q) \wedge \sim q)(1 + B(\sim p)) \\ &\equiv 1 + (1 + p + pq)(1 + q)(1 + (1 + p)) \\ &\equiv 1 + (1 + p + pq)(1 + q)(2 + p) \equiv 1 + (1 + p(1 + q))(1 + q)p \\ &\equiv 1 + ((1 + q) + p(1 + q)^2)p \equiv 1 + ((1 + q) + p(1 + q))p \\ &\equiv 1 + (1 + q)(1 + p)p \equiv 1 + (1 + q)(p + p^2) \equiv 1 + (1 + q)(p + p) \\ &\equiv 1 + (1 + q)(2p) \equiv 1 + (1 + q)(0) \equiv 1(\text{mod } 2) \blacksquare \end{aligned}$$

$$2.20) (p \rightarrow q) \wedge (r \rightarrow s) \Rightarrow (\sim q \vee \sim s \rightarrow \sim p \vee \sim r),$$

$$(p \rightarrow q) \wedge (r \rightarrow s) \Rightarrow (\sim q \wedge \sim s \rightarrow \sim p \wedge \sim r) \text{ (Destructive Dilemmas)}$$

Proof)

$$\begin{aligned}
1) & B((p \rightarrow q) \wedge (r \rightarrow s) \rightarrow (\sim q \vee \sim s \rightarrow \sim p \vee \sim r)) \\
& \equiv 1 + B((p \rightarrow q) \wedge (r \rightarrow s))(1 + B(\sim q \vee \sim s \rightarrow \sim p \vee \sim r)) \\
& \equiv 1 + (1 + p + q)(1 + r + s) \left(1 + \left(1 + B(\sim q \vee \sim s)(1 + B(\sim p \vee \sim r)) \right) \right) \\
& \equiv 1 + (1 + p + q)(1 + r + s) \left(1 + \left(1 + (1 + qs)(1 + (1 + pr)) \right) \right) \\
& \quad (\because \forall x, y: \text{proposition}, B(\sim x \vee \sim y) \equiv B(\sim x) + B(\sim y) + B(\sim x)B(\sim y)) \\
& \quad \equiv (1 + x) + (1 + y) + (1 + x)(1 + y) \equiv 2 + x + y + (1 + x + y + xy) \\
& \quad \equiv 1 + 2x + 2y + xy \equiv 1 + xy \pmod{2}) \\
& \equiv 1 + (1 + p + q)(1 + r + s)(1 + qs)pr \\
& \equiv 1 + ((1 + p + q + r + pr + qr + s + ps + qs) \\
& \quad + (qs + pqs + q^2s + rqs + prqs + q^2rs + qs^2 + ps^2q + q^2s^2))pr \\
& \equiv 1 + ((1 + p + q + r + pr + qr + s + ps + qs) \\
& \quad + (qs + pqs + qs + rqs + prqs + qrs + qs + psq + qs))pr \\
& \equiv 1 + (1 + p + q + r + pr + qr + s + ps + 5qs + 2pqs + 2rqs + prqs)pr \\
& \equiv 1 + (1 + p + q + r + pr + qr + s + ps + qs + prqs)pr \\
& \equiv 1 + (pr + p^2r + qpr + r^2p + p^2r^2 + qpr^2 + spr + p^2sr + qspr + p^2r^2qs) \\
& \equiv 1 + (pr + pr + qpr + rp + pr + qpr + spr + psr + qspr + prqs) \\
& \equiv 1 + (2pr + 2qpr + 2pr + 2spr + 2qspr) \equiv 1 \pmod{2} \blacksquare
\end{aligned}$$

$$\begin{aligned}
2) & B((p \rightarrow q) \wedge (r \rightarrow s) \rightarrow (\sim q \wedge \sim s \rightarrow \sim p \wedge \sim r)) \\
& \equiv 1 + B((p \rightarrow q) \wedge (r \rightarrow s))(1 + B(\sim q \wedge \sim s \rightarrow \sim p \wedge \sim r)) \\
& \equiv 1 + (1 + p + q)(1 + r + s) \left(1 + \left(1 + (1 + q)(1 + s)(1 + (1 + p)(1 + r)) \right) \right) \\
& \equiv 1 + (1 + p + q)(1 + r + s) \left(2 + (1 + q)(1 + s)(1 + (1 + p)(1 + r)) \right) \\
& \equiv 1 + (1 + p + q)(1 + r + s)(1 + q)(1 + s)(1 + (1 + p)(1 + r)) \\
& \equiv 1 + (p + (1 + q))(1 + q)(r + (1 + s))(1 + s)(1 + (1 + p)(1 + r)) \\
& \equiv 1 + (p(1 + q) + (1 + q)^2)(r(1 + s) + (1 + s)^2)(1 + (1 + p)(1 + r)) \\
& \equiv 1 + (p(1 + q) + (1 + q))(r(1 + s) + (1 + s))(1 + (1 + p)(1 + r)) \\
& \equiv 1 + (1 + p)(1 + q)(1 + r)(1 + s)(1 + (1 + p)(1 + r)) \\
& \equiv 1 + (1 + p)(1 + q)(1 + r)(1 + s) + (1 + p)^2(1 + q)(1 + r)^2(1 + s) \\
& \equiv 1 + (1 + p)(1 + q)(1 + r)(1 + s) + (1 + p)(1 + q)(1 + r)(1 + s) \\
& \equiv 1 + 2(1 + p)(1 + q)(1 + r)(1 + s) \equiv 1 \pmod{2} \blacksquare
\end{aligned}$$

$$2.21) (p \rightarrow q) \wedge (r \rightarrow s) \Rightarrow (p \vee r \rightarrow q \vee s),$$

$$(p \rightarrow q) \wedge (r \rightarrow s) \Rightarrow (p \wedge r \rightarrow q \wedge s) \text{ (Constructive Dilemmas)}$$

Proof)

$$\begin{aligned} 1) & B((p \rightarrow q) \wedge (r \rightarrow s) \rightarrow (p \vee r \rightarrow q \vee s)) \\ & \equiv 1 + B((p \rightarrow q) \wedge (r \rightarrow s))(1 + B(p \vee r \rightarrow q \vee s)) \\ & \equiv 1 + B(p \rightarrow q)B(r \rightarrow s) \left(1 + \left(1 + B(p \vee r)(1 + B(q \vee s)) \right) \right) \\ & \equiv 1 + B(p \rightarrow q)B(r \rightarrow s) \left(2 + B(p \vee r)(1 + B(q \vee s)) \right) \\ & \equiv 1 + B(p \rightarrow q)B(r \rightarrow s)B(p \vee r)(1 + B(q \vee s)) \\ & \equiv 1 + (1 + p + pq)(1 + r + rs)(p + r + pr)(1 + (q + s + qs)) \\ & \equiv 1 + (1 + p + pq)(1 + r + rs)(p + r + pr)(1 + q)(1 + s) \\ & \equiv 1 + (1 + p(1 + q))(1 + q)(1 + r(1 + s))(1 + s)(p + r + pr) \\ & \equiv 1 + ((1 + q) + p(1 + q)^2)((1 + s) + r(1 + s)^2)(p + r + pr) \\ & \equiv 1 + ((1 + q) + p(1 + q))((1 + s) + r(1 + s))(p + r + pr) \\ & \equiv 1 + (1 + q)(1 + p)(1 + s)(1 + r)(1 + 1 + p + r + pr) \\ & \equiv 1 + (1 + q)(1 + p)(1 + s)(1 + r)(1 + (1 + p)(1 + r)) \\ & \equiv 1 + (1 + q)(1 + p)(1 + s)(1 + r) + (1 + q)(1 + p)^2(1 + s)(1 + r)^2 \\ & \equiv 1 + (1 + q)(1 + p)(1 + s)(1 + r) + (1 + q)(1 + p)(1 + s)(1 + r) \\ & \equiv 1 + 2(1 + q)(1 + p)(1 + s)(1 + r) \equiv 1 \pmod{2} \blacksquare \end{aligned}$$

$$\begin{aligned} 2) & B((p \rightarrow q) \wedge (r \rightarrow s) \rightarrow (p \wedge r \rightarrow q \wedge s)) \\ & \equiv 1 + B((p \rightarrow q) \wedge (r \rightarrow s))(1 + B(p \wedge r \rightarrow q \wedge s)) \\ & \equiv 1 + B(p \rightarrow q)B(r \rightarrow s) \left(1 + \left(1 + B(p \wedge r)(1 + B(q \wedge s)) \right) \right) \\ & \equiv 1 + B(p \rightarrow q)B(r \rightarrow s) \left(2 + B(p \wedge r)(1 + B(q \wedge s)) \right) \\ & \equiv 1 + B(p \rightarrow q)B(r \rightarrow s)B(p \wedge r)(1 + B(q \wedge s)) \\ & \equiv 1 + (1 + p + pq)(1 + r + rs)(pr)(1 + qs) \\ & \equiv 1 + (1 + r + rs + p + pr + prs + pq + pqr + pqrs)(pr)(1 + qs) \\ & \equiv 1 + (pr + pr^2 + r^2sp + p^2r + p^2r^2 + p^2r^2s + p^2qr + p^2qr^2 + p^2qr^2s)(1 + qs) \\ & \equiv 1 + (pr + pr + rsp + pr + pr + prs + pqr + pqr + pqrs)(1 + qs) \\ & \equiv 1 + (4pr + 2rsp + 2pqr + pqrs)(1 + qs) \\ & \equiv 1 + (pqrs)(1 + qs) \equiv 1 + pqrs + pq^2rs^2 \equiv 1 + pqrs + pqrs \equiv 1 + 2pqrs \\ & \equiv 1 \pmod{2} \blacksquare \end{aligned}$$

3. Solving Boolean equation

Lemma 3.1) $\forall n, i \in \mathbb{N}, \forall x_i$: proposition

$$\begin{aligned} \left\{ (x_1, x_2, \dots, x_n) \mid \prod_{i=1}^n B(x_i) \equiv 1(\text{mod } 2) \right\} &= \left(\bigcap_{i=1}^n \{ (x_i) \mid B(x_i) \equiv 1(\text{mod } 2) \} \right) \\ &= \{ (x_1, x_2, \dots, x_n) \mid (1, 1, 1, \dots, 1) \} \\ \left\{ (x_1, x_2, \dots, x_n) \mid \prod_{i=1}^n B(x_i) \equiv 0(\text{mod } 2) \right\} &= \left(\bigcap_{i=1}^n \{ (x_i) \mid B(x_i) \equiv 1(\text{mod } 2) \} \right)^c \\ &= \{ (x_1, x_2, \dots, x_n) \mid (1, 1, 1, \dots, 1) \}^c \end{aligned}$$

Proof)

Let $\forall n, i \in \mathbb{N}, \forall x_i$: proposition

For $1 \leq i \leq n, \{ (x_i) \mid (0), (1) \} = \{ x_i \mid 0, 1 \}$

Assume $\exists B(x_i) \equiv 0(\text{mod } 2) \Rightarrow \prod_{i=1}^n B(x_i) \equiv B(x_1)B(x_2) \dots B(x_n) \equiv 0(\text{mod } 2)$

Therefore $\prod_{i=1}^n B(x_i) \equiv 1(\text{mod } 2) \Rightarrow \forall B(x_i) \equiv 1(\text{mod } 2)$

$$\begin{aligned} \Rightarrow \left\{ (x_1, x_2, \dots, x_n) \mid \prod_{i=1}^n B(x_i) \equiv 1(\text{mod } 2) \right\} &= \left(\bigcap_{i=1}^n \{ (x_i) \mid B(x_i) \equiv 1(\text{mod } 2) \} \right) \\ &= \{ (x_1, x_2, \dots, x_n) \mid (1, 1, 1, \dots, 1) \} \end{aligned}$$

And $\prod_{i=1}^n B(x_i) \equiv 0(\text{mod } 2) \Rightarrow \sim(\forall B(x_i) \equiv 1(\text{mod } 2))$

$$\begin{aligned} \Rightarrow \left\{ (x_1, x_2, \dots, x_n) \mid \prod_{i=1}^n B(x_i) \equiv 0(\text{mod } 2) \right\} &= \left(\bigcap_{i=1}^n \{ (x_i) \mid B(x_i) \equiv 1(\text{mod } 2) \} \right)^c \\ &= \{ (x_1, x_2, \dots, x_n) \mid (1, 1, 1, \dots, 1) \}^c \blacksquare \end{aligned}$$

Method) Steps of solving Boolean equation

Step 1) We change Boolean equation to product of mod 2 using by 1.*.

Step 2) We get sets of solution using by Lemma3.1.

Example) $\forall n, i \in \mathbb{N}, \forall x_i$: proposition

- *Ex 1)* $B(x_1 \wedge \sim x_2 \wedge x_3) \equiv 1(mod 2)$

Step 1)

$$B(x_1 \wedge \sim x_2 \wedge x_3) \equiv 1(mod 2) \Rightarrow x_1(1 + x_2)x_3 \equiv 1(mod 2)$$

Step 2)

$$\begin{aligned} \Rightarrow \left(\begin{array}{l} \{(x_1) \mid x_1 \equiv 1(mod 2)\} \\ \cap \{(x_2) \mid 1 + x_2 \equiv 1(mod 2)\} \\ \cap \{(x_3) \mid x_3 \equiv 1(mod 2)\} \end{array} \right) &= \left(\begin{array}{l} \{(x_1) \mid x_1 \equiv 1(mod 2)\} \\ \cap \{(x_2) \mid x_2 \equiv 0(mod 2)\} \\ \cap \{(x_3) \mid x_3 \equiv 1(mod 2)\} \end{array} \right) \\ &= \{(x_1, x_2, x_3) \mid (1,0,1)\} \end{aligned}$$

- *Ex 2)* $B(x_1 \vee x_2 \vee x_3) \equiv 1(mod 2)$

Step 1)

$$B(x_1 \vee x_2 \vee x_3) \equiv 1(mod 2) \Rightarrow 1 + (1 + x_1)(1 + x_2)(1 + x_3) \equiv 1(mod 2)$$

$$\Rightarrow (1 + x_1)(1 + x_2)(1 + x_3) \equiv 0(mod 2)$$

Step 2)

$$\begin{aligned} \Rightarrow \left(\begin{array}{l} \{(x_1) \mid 1 + x_1 \equiv 1(mod 2)\} \\ \cap \{(x_2) \mid 1 + x_2 \equiv 1(mod 2)\} \\ \cap \{(x_3) \mid 1 + x_3 \equiv 1(mod 2)\} \end{array} \right)^c &= \left(\begin{array}{l} \{(x_1) \mid x_1 \equiv 0(mod 2)\} \\ \cap \{(x_2) \mid x_2 \equiv 0(mod 2)\} \\ \cap \{(x_3) \mid x_3 \equiv 0(mod 2)\} \end{array} \right)^c \\ &= \{(x_1, x_2, x_3) \mid (0,0,0)\}^c \\ &= \{(x_1, x_2, x_3) \mid (0,0,1), (0,1,0), (0,1,1), (1,0,0), (1,0,1), (1,1,0), (1,1,1)\} \end{aligned}$$

- Ex 3) $B(x_1 \vee (\sim x_2 \wedge x_3)) \equiv 1(\text{mod } 2)$

Step 1)

$$\begin{aligned} B(x_1 \vee (\sim x_2 \wedge x_3)) &\equiv 1(\text{mod } 2) \Rightarrow 1 + (1 + x_1)(1 + B(\sim x_2 \wedge x_3)) \equiv 1(\text{mod } 2) \\ &\Rightarrow 1 + (1 + x_1)(1 + (1 + x_2)x_3) \equiv 1(\text{mod } 2) \Rightarrow (1 + x_1)(1 + (1 + x_2)x_3) \equiv 0(\text{mod } 2) \end{aligned}$$

Step 2)

$$\begin{aligned} &\Rightarrow \left(\begin{array}{l} \{(x_1) \mid 1 + x_1 \equiv 1(\text{mod } 2)\} \\ \cap \{(x_2, x_3) \mid 1 + (1 + x_2)x_3 \equiv 1(\text{mod } 2)\} \end{array} \right)^c \\ &= \left(\begin{array}{l} \{(x_1) \mid x_1 \equiv 0(\text{mod } 2)\} \\ \cap \{(x_2, x_3) \mid (1 + x_2)x_3 \equiv 0(\text{mod } 2)\} \end{array} \right)^c \\ &= \left(\begin{array}{l} \{(x_1) \mid x_1 \equiv 0(\text{mod } 2)\} \\ \cap \left(\begin{array}{l} \{(x_2) \mid 1 + x_2 \equiv 1(\text{mod } 2)\} \\ \cap \{(x_3) \mid x_3 \equiv 1(\text{mod } 2)\} \end{array} \right)^c \end{array} \right)^c = \left(\begin{array}{l} \{(x_1) \mid x_1 \equiv 0(\text{mod } 2)\} \\ \cap \left(\begin{array}{l} \{(x_2) \mid x_2 \equiv 0(\text{mod } 2)\} \\ \cap \{(x_3) \mid x_3 \equiv 1(\text{mod } 2)\} \end{array} \right)^c \end{array} \right)^c \\ &= \left(\begin{array}{l} \{(x_1) \mid (0)\} \\ \cap \{(x_2, x_3) \mid (0,1)\}^c \end{array} \right)^c = \left(\begin{array}{l} \{(x_1) \mid (0)\} \\ \cap \{(x_2, x_3) \mid (0,0), (1,0), (1,1)\} \end{array} \right)^c \\ &= \{(x_1, x_2, x_3) \mid (0,0,0), (0,1,0), (0,1,1)\}^c \\ &= \{(x_1, x_2, x_3) \mid (0,0,1), (1,0,0), (1,0,1), (1,1,0), (1,1,1)\} \end{aligned}$$

- Ex 4) $B((x_1 \wedge x_2 \wedge \sim x_3) \vee (\sim x_1 \wedge x_2 \wedge \sim x_3)) \equiv 1(\text{mod } 2)$

Step 1)

$$\begin{aligned} B((x_1 \wedge x_2 \wedge \sim x_3) \vee (\sim x_1 \wedge x_2 \wedge \sim x_3)) &\equiv 1(\text{mod } 2) \\ &\Rightarrow 1 + (1 + B(x_1 \wedge x_2 \wedge \sim x_3))(1 + B(\sim x_1 \wedge x_2 \wedge \sim x_3)) \equiv 1(\text{mod } 2) \\ &\Rightarrow 1 + (1 + x_1x_2(1 + x_3))(1 + (1 + x_1)x_2(1 + x_3)) \equiv 1(\text{mod } 2) \\ &\Rightarrow (1 + x_1x_2(1 + x_3))(1 + (1 + x_1)x_2(1 + x_3)) \equiv 0(\text{mod } 2) \end{aligned}$$

Step 2)

$$\begin{aligned} &\Rightarrow \left(\begin{array}{l} \{(x_1, x_2, x_3) \mid 1 + x_1x_2(1 + x_3) \equiv 1(\text{mod } 2)\} \\ \cap \{(x_1, x_2, x_3) \mid 1 + (1 + x_1)x_2(1 + x_3) \equiv 1(\text{mod } 2)\} \end{array} \right)^c \\ &= \left(\begin{array}{l} \{(x_1, x_2, x_3) \mid x_1x_2(1 + x_3) \equiv 0(\text{mod } 2)\} \\ \cap \{(x_1, x_2, x_3) \mid (1 + x_1)x_2(1 + x_3) \equiv 0(\text{mod } 2)\} \end{array} \right)^c \\ &= \left(\begin{array}{l} \left(\begin{array}{l} \{(x_1) \mid x_1 \equiv 1(\text{mod } 2)\} \\ \cap \{(x_2) \mid x_2 \equiv 1(\text{mod } 2)\} \\ \cap \{(x_3) \mid (1 + x_3) \equiv 1(\text{mod } 2)\} \end{array} \right)^c \\ \cap \left(\begin{array}{l} \{(x_1) \mid (1 + x_1) \equiv 1(\text{mod } 2)\} \\ \cap \{(x_2) \mid x_2 \equiv 1(\text{mod } 2)\} \\ \cap \{(x_3) \mid (1 + x_3) \equiv 1(\text{mod } 2)\} \end{array} \right)^c \end{array} \right)^c = \left(\begin{array}{l} \left(\begin{array}{l} \{(x_1) \mid x_1 \equiv 1(\text{mod } 2)\} \\ \cap \{(x_2) \mid x_2 \equiv 1(\text{mod } 2)\} \\ \cap \{(x_3) \mid x_3 \equiv 0(\text{mod } 2)\} \end{array} \right)^c \\ \cap \left(\begin{array}{l} \{(x_1) \mid x_1 \equiv 0(\text{mod } 2)\} \\ \cap \{(x_2) \mid x_2 \equiv 1(\text{mod } 2)\} \\ \cap \{(x_3) \mid x_3 \equiv 0(\text{mod } 2)\} \end{array} \right)^c \end{array} \right)^c \\ &= \left(\begin{array}{l} \{(x_1, x_2, x_3) \mid (1,1,0)\}^c \\ \cap \{(x_1, x_2, x_3) \mid (0,1,0)\}^c \end{array} \right)^c = \left(\begin{array}{l} \{(x_1, x_2, x_3) \mid (1,1,0)\} \\ \cup \{(x_1, x_2, x_3) \mid (0,1,0)\} \end{array} \right) \\ &= \{(x_1, x_2, x_3) \mid (0,1,0), (1,1,0)\} \end{aligned}$$

4. Application

4.1. Proof of validity

Example) H_1, H_2 : Hypothesis, C : Conclusion

$$H_1 : x_1 \vee (x_2 \wedge x_3)$$

$$H_2 : x_1 \vee x_2 \rightarrow x_4$$

$$C : x_1 \vee x_4$$

$$H_1 \wedge H_2 \Rightarrow C \Leftrightarrow B(H_1 \wedge H_2 \rightarrow C) \equiv 1(\text{mod } 2)$$

$$\Rightarrow 1 + B(x_1 \vee (x_2 \wedge x_3))B(x_1 \vee x_2 \rightarrow x_4)(1 + B(x_1 \vee x_4)) \equiv 1(\text{mod } 2)$$

$$B(x_1 \vee (x_2 \wedge x_3))B(x_1 \vee x_2 \rightarrow x_4)(1 + B(x_1 \vee x_4)) \equiv 0(\text{mod } 2) \Rightarrow$$

$$(1 + (1 + x_1)(1 + x_2x_3))(1 + (1 + (1 + x_1)(1 + x_2))(1 + x_4))(1 + (1 + (1 + x_1)(1 + x_4)))$$

$$\equiv (1 + (1 + x_1)(1 + x_2x_3))(1 + (1 + (1 + x_1)(1 + x_2))(1 + x_4))((1 + x_1)(1 + x_4))$$

$$\equiv 0(\text{mod } 2)$$

- Method 1) deployment

Let $1 + p \equiv p'$ for $\forall p$: proposition

$$(1 + (1 + x_1)(1 + x_2x_3))(1 + (1 + (1 + x_1)(1 + x_2))(1 + x_4))((1 + x_1)(1 + x_4))$$

$$\equiv (1 + x_1'(x_2x_3)')(1 + (1 + x_1'x_2')x_4')(x_1'x_4')$$

$$\equiv (1 + x_1'(x_2x_3)')(1 + x_4' + x_1'x_2'x_4')(x_1'x_4')$$

$$\equiv (1 + x_1'(x_2x_3)')(x_1'x_4' + x_1'x_4'x_4' + x_1'x_4'x_1'x_2'x_4')$$

$$\equiv (1 + x_1'(x_2x_3)')(x_1'x_4' + x_1'x_4' + x_1'x_2'x_4') \equiv (1 + x_1'(x_2x_3)')(2x_1'x_4' + x_1'x_2'x_4')$$

$$\equiv (1 + x_1'(x_2x_3)')(x_1'x_2'x_4') \equiv x_1'x_2'x_4' + x_1'x_2'x_4'x_1'(x_2x_3)'$$

$$\equiv x_1'x_2'x_4' + x_1'x_2'x_4'(x_2x_3)' \equiv x_1'x_2'x_4' + x_1'x_2'x_4'(1 + x_2x_3)$$

$$\equiv x_1'x_2'x_4' + x_1'x_2'x_4' + x_1'x_2'x_4'x_2x_3 \equiv 2x_1'x_2'x_4' + x_1'x_2'x_4'x_2x_3$$

$$\equiv x_1'(1 + x_2)x_4'x_2x_3 \equiv x_1'(x_2 + x_2^2)x_4'x_3 \equiv x_1'(x_2 + x_2)x_4'x_3$$

$$\equiv x_1'(2x_2)x_4'x_3 \equiv x_1'(0)x_4'x_3 \equiv 0(\text{mod } 2)$$

$$\therefore B(H_1 \wedge H_2 \rightarrow C) \equiv 1(\text{mod } 2)$$

- Method 2) using by solving method of 3.

$$\begin{aligned}
& (1 + (1 + x_1)(1 + x_2x_3)) \left(1 + (1 + (1 + x_1)(1 + x_2))(1 + x_4)\right) \left((1 + x_1)(1 + x_4)\right) \\
& \quad \equiv 0(\text{mod } 2) \\
& \Rightarrow \left(\begin{array}{l} \{(x_1, x_2, x_3) \mid 1 + (1 + x_1)(1 + x_2x_3) \equiv 1(\text{mod } 2)\} \\ \cap \{(x_1, x_2, x_4) \mid 1 + (1 + (1 + x_1)(1 + x_2))(1 + x_4) \equiv 1(\text{mod } 2)\} \\ \cap \{(x_1, x_4) \mid (1 + x_1)(1 + x_4) \equiv 1(\text{mod } 2)\} \end{array} \right)^c \\
& = \left(\begin{array}{l} \{(x_1, x_2, x_3) \mid (1 + x_1)(1 + x_2x_3) \equiv 0(\text{mod } 2)\} \\ \cap \{(x_1, x_2, x_4) \mid (1 + (1 + x_1)(1 + x_2))(1 + x_4) \equiv 0(\text{mod } 2)\} \\ \cap \{(x_1, x_4) \mid (1 + x_1)(1 + x_4) \equiv 1(\text{mod } 2)\} \end{array} \right)^c \\
& = \left(\begin{array}{l} \left(\begin{array}{l} \{(x_1) \mid (1 + x_1) \equiv 1(\text{mod } 2)\} \\ \cap \{(x_2, x_3) \mid (1 + x_2x_3) \equiv 1(\text{mod } 2)\} \end{array} \right)^c \\ \cap \left(\begin{array}{l} \{(x_1, x_2) \mid 1 + (1 + x_1)(1 + x_2) \equiv 1(\text{mod } 2)\} \\ \cap \{(x_4) \mid 1 + x_4 \equiv 1(\text{mod } 2)\} \end{array} \right)^c \\ \cap \left(\begin{array}{l} \{(x_1) \mid (1 + x_1) \equiv 1(\text{mod } 2)\} \\ \cap \{(x_4) \mid (1 + x_4) \equiv 1(\text{mod } 2)\} \end{array} \right) \end{array} \right)^c \\
& = \left(\begin{array}{l} \left(\begin{array}{l} \{(x_1) \mid x_1 \equiv 0(\text{mod } 2)\} \\ \cap \{(x_2, x_3) \mid x_2x_3 \equiv 0(\text{mod } 2)\} \end{array} \right)^c \\ \cap \left(\begin{array}{l} \{(x_1, x_2) \mid (1 + x_1)(1 + x_2) \equiv 0(\text{mod } 2)\} \\ \cap \{(x_4) \mid x_4 \equiv 0(\text{mod } 2)\} \end{array} \right)^c \\ \cap \left(\begin{array}{l} \{(x_1) \mid x_1 \equiv 0(\text{mod } 2)\} \\ \cap \{(x_4) \mid x_4 \equiv 0(\text{mod } 2)\} \end{array} \right) \end{array} \right)^c \\
& = \left(\begin{array}{l} \left(\begin{array}{l} \{(x_1) \mid x_1 \equiv 0(\text{mod } 2)\} \\ \cap \left(\begin{array}{l} \{(x_2) \mid x_2 \equiv 1(\text{mod } 2)\} \\ \cap \{(x_3) \mid x_3 \equiv 1(\text{mod } 2)\} \end{array} \right)^c \end{array} \right)^c \\ \cap \left(\begin{array}{l} \left(\begin{array}{l} \{(x_1) \mid (1 + x_1) \equiv 1(\text{mod } 2)\} \\ \cap \{(x_2) \mid (1 + x_2) \equiv 1(\text{mod } 2)\} \end{array} \right)^c \\ \cap \{(x_4) \mid x_4 \equiv 0(\text{mod } 2)\} \end{array} \right)^c \\ \cap \left(\begin{array}{l} \{(x_1) \mid x_1 \equiv 0(\text{mod } 2)\} \\ \cap \{(x_4) \mid x_4 \equiv 0(\text{mod } 2)\} \end{array} \right) \end{array} \right)^c \\
& = \left(\begin{array}{l} \left(\begin{array}{l} \{(x_1) \mid x_1 \equiv 0(\text{mod } 2)\} \\ \cap \left(\begin{array}{l} \{(x_2) \mid x_2 \equiv 1(\text{mod } 2)\} \\ \cap \{(x_3) \mid x_3 \equiv 1(\text{mod } 2)\} \end{array} \right)^c \end{array} \right)^c \\ \cap \left(\begin{array}{l} \left(\begin{array}{l} \{(x_1) \mid x_1 \equiv 0(\text{mod } 2)\} \\ \cap \{(x_2) \mid x_2 \equiv 0(\text{mod } 2)\} \end{array} \right)^c \\ \cap \{(x_4) \mid x_4 \equiv 0(\text{mod } 2)\} \end{array} \right)^c \\ \cap \left(\begin{array}{l} \{(x_1) \mid x_1 \equiv 0(\text{mod } 2)\} \\ \cap \{(x_4) \mid x_4 \equiv 0(\text{mod } 2)\} \end{array} \right) \end{array} \right)^c = \left(\begin{array}{l} \left(\begin{array}{l} \{(x_1) \mid (0)\} \\ \cap \left(\begin{array}{l} \{(x_2) \mid (1)\} \\ \cap \{(x_3) \mid (1)\} \end{array} \right)^c \end{array} \right)^c \\ \cap \left(\begin{array}{l} \left(\begin{array}{l} \{(x_1) \mid (0)\} \\ \cap \{(x_2) \mid (0)\} \end{array} \right)^c \\ \cap \{(x_4) \mid (0)\} \end{array} \right)^c \\ \cap \left(\begin{array}{l} \{(x_1) \mid (0)\} \\ \cap \{(x_4) \mid (0)\} \end{array} \right) \end{array} \right)^c
\end{aligned}$$

$$\begin{aligned}
&= \left(\left(\begin{array}{l} \{(x_1) | (0)\} \\ \cap \{(x_2, x_3) | (1,1)\}^c \end{array} \right)^c \right)^c = \left(\begin{array}{l} \{(x_1) | (0)\} \\ \cap \{(x_2, x_3) | (0,0), (0,1), (1,0)\} \end{array} \right)^c \right)^c \\
&= \left(\begin{array}{l} \{(x_1, x_2) | (0,0)\}^c \\ \cap \{(x_4) | (0)\} \end{array} \right)^c \cap \{(x_1, x_2) | (0,1), (1,0), (1,1)\}^c \\
&= \left(\begin{array}{l} \{(x_1, x_2, x_3) | (0,0,0), (0,0,1), (0,1,0)\}^c \\ \cap \{(x_1, x_2, x_4) | (0,1,0), (1,0,0), (1,1,0)\}^c \\ \cap \{(x_1, x_4) | (0,0)\} \end{array} \right)^c \\
&= \left(\begin{array}{l} \{(x_1, x_2, x_3) | (0,1,1), (1,0,0), (1,0,1), (1,1,0), (1,1,1)\} \\ \cap \{(x_1, x_2, x_4) | (0,0,0), (0,0,1), (0,1,1), (1,0,1), (1,1,1)\} \\ \cap \{(x_1, x_4) | (0,0)\} \end{array} \right)^c \\
&= \left(\begin{array}{l} \{(x_1, x_2, x_3, x_4) | (0,1,1,1), (1,0,0,1), (1,0,1,1), (1,1,0,1), (1,1,1,1)\} \\ \cap \{(x_1, x_4) | (0,0)\} \end{array} \right)^c \\
&= \emptyset^c = U
\end{aligned}$$

Let $\{(x)|(*)\} = \{(x)|(0), (1)\} \Rightarrow U = \{(x_1, x_2, x_3, x_4) | (*,*,*,*)\}$

$\therefore \{(x_1, x_2, x_3, x_4) | B(H_1 \wedge H_2 \rightarrow C) \equiv 1(\text{mod } 2)\} = \{(x_1, x_2, x_3, x_4) | (*,*,*,*)\}$

4.2. Simplification

Example)

x	y	z	$F(x, y, z)$
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	0

$$B\left(\left((\sim x) \wedge y \wedge (\sim z)\right) \vee \left(x \wedge (\sim y) \wedge (\sim z)\right) \vee \left(x \wedge y \wedge (\sim z)\right)\right) \text{ (by SOP(Sum of Product))}$$

$$\begin{aligned} &\equiv 1 + \left(1 + B((\sim x) \wedge y \wedge (\sim z))\right) \left(1 + B(x \wedge (\sim y) \wedge (\sim z))\right) \left(1 + B(x \wedge y \wedge (\sim z))\right) \\ &\equiv 1 + \left(1 + (1 + x)y(1 + z)\right) \left(1 + x(1 + y)(1 + z)\right) \left(1 + xy(1 + z)\right) \\ &\equiv 1 + \left(1 + (1 + x)y(1 + z) + x(1 + y)(1 + z) + (1 + x)y(1 + z)x(1 + y)(1 + z)\right) \left(1 + xy(1 + z)\right) \\ &\equiv 1 + \left(1 + (1 + x)y(1 + z) + x(1 + y)(1 + z) + (x + x^2)(y + y^2)(1 + z)^2\right) \left(1 + xy(1 + z)\right) \\ &\equiv 1 + \left(1 + (1 + x)y(1 + z) + x(1 + y)(1 + z) + (x + x)(y + y)(1 + z)\right) \left(1 + xy(1 + z)\right) \\ &\equiv 1 + \left(1 + (1 + x)y(1 + z) + x(1 + y)(1 + z) + (2x)(2y)(1 + z)\right) \left(1 + xy(1 + z)\right) \\ &\equiv 1 + \left(1 + (1 + x)y(1 + z) + x(1 + y)(1 + z)\right) \left(1 + xy(1 + z)\right) \\ &\equiv 1 + \left(1 + ((1 + x)y + x(1 + y))(1 + z)\right) \left(1 + xy(1 + z)\right) \\ &\equiv 1 + \left(1 + (y + xy + x + xy)(1 + z)\right) \left(1 + xy(1 + z)\right) \\ &\equiv 1 + \left(1 + (y + x + 2xy)(1 + z)\right) \left(1 + xy(1 + z)\right) \\ &\equiv 1 + \left(1 + (y + x)(1 + z)\right) \left(1 + xy(1 + z)\right) \\ &\equiv 1 + \left(1 + (y + x)(1 + z) + xy(1 + z) + xy(y + x)(1 + z)^2\right) \\ &\equiv 1 + \left(1 + (y + x)(1 + z) + xy(1 + z) + xy(y + x)(1 + z)\right) \\ &\equiv 1 + \left(1 + ((y + x) + xy + xy(y + x))(1 + z)\right) \\ &\equiv 1 + \left(1 + (y + x + xy + xy^2 + x^2y)(1 + z)\right) \\ &\equiv 1 + \left(1 + (y + x + xy + xy + xy)(1 + z)\right) \\ &\equiv 1 + \left(1 + (y + x + xy)(1 + z)\right) \equiv 1 + 1 + (y + x + xy)(1 + z) \\ &\equiv (y + x + xy)(1 + z) \pmod{2} \\ &\therefore F(x, y, z) \equiv (y + x + xy)(1 + z) \equiv B((x \vee y) \wedge (\sim z)) \pmod{2} \end{aligned}$$

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