

Refutation of definable modal operators on stable set lattices

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Abstract: We evaluate definable modal operators on stable set lattices. None of the definitions or implementations is tautologous, hence refuting the operators.

We assume the method and apparatus of Meth8/VL4 with Tautology as the designated *proof* value, **F** as contradiction, **N** as truthity (non-contingency), and **C** as falsity (contingency). The 16-valued truth table is row-major and horizontal. The 16-valued truth table is row-major and horizontal, or repeating fragments of 128-tables, sometimes with table counts, for more variables. (See ersatz-systems.com.)

LET \sim Not, \neg ; $+$ Or, \vee ; $-$ Not Or; $\&$ And, \wedge ; \setminus Not And;
 $>$ Imply, greater than, \rightarrow ; $<$ Not Imply, less than, \Leftarrow
 $=$ Equivalent, \equiv ; $@$ Not Equivalent, \neq ;
 $\%$ possibility, for one or some, \exists , \diamond ; $\#$ necessity, for every or all, \forall , \square ;
 $\sim(y < x)$ ($x \leq y$), ($x \subseteq y$).

From: Goldblatt, R. (2018). Definable operators on stable set lattices.
 arxiv.org/pdf/1812.01264.pdf rob.goldblatt@msor.vuw.ac.nz

The key idea is that of a first-order definable operation on a stable set lattice, an idea that goes to the heart of Kripke's semantical interpretation of the modalities \square and \diamond . On the algebra of subsets of a Kripke frame (X, R) , the modal connectives can be interpreted as operations assigning to each set $A \subseteq X$ the sets

[LET $p, q, r, s: A, R, x, y.$]

$$\square A = \{x : \forall y(xRy \rightarrow y \in A)\} \quad (2.1.1)$$

$$\#p = (r > ((r \& (q \& \#s)) > (r < p))) ; \mathbf{FNFN} \ \mathbf{FNFN} \ \mathbf{FNFN} \ \mathbf{FNFN} \quad (2.1.2)$$

$$\text{and } \diamond A = \{x : \exists y(xRy \ \& \ y \in A)\}. \quad (2.2.1)$$

$$\%p = (r > ((r \& (q \& \%s)) > (r < p))) ; \quad \mathbf{CTCT} \ \mathbf{CTCN} \ \mathbf{CTCT} \ \mathbf{CTCF} \quad (2.2.2)$$

Example 3.2. Modal operators.

[LET $p, r, t, x, y, z: A, R, T, x, y, z.$]

$$\square A = \{x : \forall y[\forall z(z \in A \rightarrow zRy) \rightarrow xTy]\} \quad (3.2.1.1)$$

$$\#p = (x > (((\#z < p) > (\#z \& (r \& \#y))) > (x \& (t \& \#y)))) ; \quad (3.2.1.2)$$

$\mathbf{TCTC} \ \mathbf{TCTC} \ \mathbf{TCTC} \ \mathbf{TCTC}, \mathbf{CCCC} \ \mathbf{CCCC} \ \mathbf{CCCC} \ \mathbf{CCCC},$
 $\mathbf{FNFN} \ \mathbf{FNFN} \ \mathbf{FNFN} \ \mathbf{FNFN}, \mathbf{CCCC} \ \mathbf{TCTC} \ \mathbf{CCCC} \ \mathbf{TCTC},$
 $\mathbf{CTCT} \ \mathbf{CTCT} \ \mathbf{CTCT} \ \mathbf{CTCT}$

$$\Diamond A = \{x : \forall y [\forall z (z \in A \rightarrow zTy) \rightarrow xRy] \}, \quad (3.2.2.1)$$

$$\begin{aligned} \%p = & (x > (((\#z < p) > (\#z \& (t \& \#y))) > (x \& (r \& \#y)))) ; \\ & \text{CTCT CTCT CTCT CTCT, } \mathbf{FFFF} \mathbf{FNFN} \mathbf{FFFF} \mathbf{FNFN}, \\ & \mathbf{NFNF} \mathbf{FNFN} \mathbf{NFNF} \mathbf{FNFN} \end{aligned} \quad (3.2.2.2)$$

Eqs. 2 and 3 as rendered are *not* tautologous, hence refuting definable modal operators on stable set lattices.