

Refutation of dynamic modal logic

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Abstract: We evaluate two definitions of dynamic modal logic on which axioms are built. The definitions are *not* tautologous and moreover logically equivalent. This refutes dynamic modal logic.

We assume the method and apparatus of Meth8/VL4 with Tautology as the designated *proof* value, **F** as contradiction, **N** as truthity (non-contingency), and **C** as falsity (contingency). The 16-valued truth table is row-major and horizontal. The 16-valued truth table is row-major and horizontal, or repeating fragments of 128-tables, sometimes with table counts, for more variables. (See ersatz-systems.com.)

LET $p, q, r, s: p, q, a, b;$
 \sim Not, \neg ; $+$ Or, \vee, \cup ; $-$ Not Or; $\&$ And, \wedge, \cap ; \setminus Not And;
 $>$ Imply, greater than, $\rightarrow, \mapsto, \succ, \supset, \vdash, \models, \rightsquigarrow$; $<$ Not Imply, less than, $\in, \prec, \subset, \leftarrow$;
 $=$ Equivalent, $\equiv, :=, \iff, \leftrightarrow, \triangleq$ $@$ Not Equivalent, \neq ;
 $\%$ possibility, for one or some, \exists, \diamond, M ; $\#$ necessity, for every or all, \forall, \square, L ;
 $(z=z)$ **T** as tautology, \top , ordinal 3; $(z@z)$ **F** as contradiction, $\emptyset, \text{Null}, \perp$, zero;
 $(\%z\<\#z)$ **C** as contingency, Δ , ordinal 1; $(\%z\>\#z)$ **N** as non-contingency, ∇ , ordinal 2;
 $\sim(y < x)$ ($x \leq y$), ($x \subseteq y$); $(A=B)$ ($A \sim B$).

From: [en.wikipedia.org/wiki/Dynamic_logic_\(modal_logic\)](http://en.wikipedia.org/wiki/Dynamic_logic_(modal_logic))

Modal logic is characterized by the modal operators $\Box p$ (box p) asserting that p is necessarily the case, and $\Diamond p$ (diamond p) asserting that p is possibly the case. Dynamic logic extends this by associating to every action a the modal operators $[a]$ and $\langle a \rangle$, thereby making it a multimodal logic. The meaning of $[a]p$ is that after performing action a it is necessarily the case that p holds, that is, a must bring about p. The meaning of $\langle a \rangle p$ is that after performing a it is possible that p holds, that is, a might bring about p. These operators are related by

$$[a]p = \sim \langle a \rangle \sim p \text{ and} \tag{0.1.1}$$

We write Eq. 0.1.1 as:

"Necessarily r applies to p implies the necessity of p" is equivalent to
 "Not possibly r applies to p implies not the possibility of p".

$$((\#r \& p) \> \#p) = ((\sim \%r \& p) \> \sim \%p); \quad \text{TCTC TTTT TCTC TTTT} \tag{0.1.2}$$

$$\langle a \rangle p = \sim [a] \sim p, \tag{0.2.1}$$

We write Eq. 0.2.1 as:

"Possibly r applies to p implies the possibility of p" is equivalent to
 "Not necessarily r applies to p implies not the necessity of p".

$$((\%r \& p) \> \%p) = ((\sim \#r \& p) \> \sim \#p); \quad \text{TCTC TTTT TCTC TTTT} \tag{0.2.2}$$

analogously to the relationship between the universal (\forall) and existential (\exists) quantifiers.

Remark 0: Eqs. 0.1.2 and 0.2.2 as rendered as definitions are *not* tautologous. In fact, they are logical equivalences. This refutes the definitions of dynamic modal logic. Those definitions as analogous to the quantifiers is thus denied.

$$\text{A3. } [a \cup b]p \equiv [a]p \wedge [b]p \tag{3.1}$$

(A3 says that if doing one of a or b must bring about p, then a must bring about p and likewise for b, and conversely.)

Remark 3.1: Eq. 3.1 is an obvious tautology and hence trivial as an axiom.

The stated definitions in Eqs. 0.1.2 and 0.2.2 are *not* tautologous and also logically identical. This on its face refutes dynamic modal logic.