

Refutation of the unification type of simple symmetric modal logics

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Abstract: We evaluate the definitions of two new modal connectives as box-plus and box-minus. Neither is tautologous and both are equivalent. This refutes the unification type of simple symmetric modal logics and implies it is a *non* tautologous fragment of the universal logic $V\mathcal{L}4$.

We assume the method and apparatus of Meth8/ $V\mathcal{L}4$ with Tautology as the designated proof value, F as contradiction, N as truthity (non-contingency), and C as falsity (contingency). The 16-valued truth table is row-major and horizontal, or repeating fragments of 128-tables, sometimes with table counts, for more variables. (See ersatz-systems.com.)

LET \sim Not, \neg ; + Or, \vee, \cup ; - Not Or; & And, \wedge, \cap ; \ Not And;
 $>$ Imply, greater than, $\rightarrow, \mapsto, \succ, \supset, \vdash, \models, \twoheadrightarrow$; $<$ Not Imply, less than, $\in, \prec, \subset, \neq, \neq, \leftarrow$;
 $=$ Equivalent, $\equiv, :=, \iff, \leftrightarrow, \triangleq$ @ Not Equivalent, \neq ;
 $\%$ possibility, for one or some, \exists, \diamond, M ; # necessity, for every or all, \forall, \square, L ;
 $(z=z)$ T as tautology, \top , ordinal 3; $(z@z)$ F as contradiction, $\emptyset, \text{Null}, \perp$, zero;
 $(\%z\#z)$ C as contingency, Δ , ordinal 1; $(\%z\#z)$ N as non-contingency, ∇ , ordinal 2;
 $\sim(y < x)$ ($x \leq y$), ($x \subseteq y$); $(A=B)$ ($A \sim B$).
 Note: For clarity we usually distribute quantifiers on each variable as designated.

From: Balbiani, P.; Gencer, Ç. (2019). About the unification type of simple symmetric modal logics. arxiv.org/pdf/1902.03770.pdf philippe.balbani@irit.fr

2 Syntax.

For all parameters p , we write “ p^0 ” to mean “ $\neg p$ ” and we write “ p^1 ” to mean “ p ”.

Let \boxplus and \boxminus be the modal connectives defined as follows: (2.0)

Remark 2.0: We name the respective symbols \boxplus and \boxminus as box-plus and box-minus.

$$\boxplus\phi ::= (p^0 \wedge q^0 \rightarrow \square(p^1 \wedge q^0 \rightarrow \square(p^0 \wedge q^1 \rightarrow \square(p^0 \wedge q^0 \rightarrow \phi))))), \quad (2.1.1)$$

LET p, q, r : $p^1=p, q^1=q, p^0=\sim p, q^0=\sim q, \phi$.

$$((\sim p \& \sim q) \> \#((p \& \sim q) \> \#((\sim p \& q) \> \#((\sim p \& \sim q) \> t)))) = (p=p); \quad (2.1.2)$$

NTTT NTTT NTTT NTTT

$$\boxminus\phi ::= (p^0 \wedge q^0 \rightarrow \square(p^0 \wedge q^1 \rightarrow \square(p^1 \wedge q^0 \rightarrow \square(p^0 \wedge q^0 \rightarrow \phi))))). \quad (2.2.1)$$

$$((\sim p \& \sim q) \> \#((\sim p \& q) \> \#((p \& \sim q) \> \#((\sim p \& \sim q) \> t)))) = (p=p); \quad (2.2.2)$$

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Eqs. 2.1.2 and 2.2.2 are *not* tautologous and logically equivalent because the respective conjunctive clauses are identical. This refutes the unification type of simple symmetric modal logics.