

Honorable Sir (beautiful lady):

Because I don't know English.

Because my education level is very low. (There are some mistakes, I believe you can modify them.)

It is proved that any odd number $【2N+1=2\times P_a+P_b】$ There are no counterexamples.

吴叶唐寅: theorem (四) $(2N+1=2+P_a、\text{ or、 } 2N+1=2\times P_a+P_b)$

Abstract hypothesis, simulate arithmetic logic of the same number, judge reasoning and hypothesis contradiction

theorem: An odd number greater than 5 equals 2 plus an odd prime.

Or an odd number greater than 5 equals two identical prime numbers plus one odd prime number(Integer theory. Contradiction)

$$(2N+1=2+P_a、\text{ or、 } 2N+1=2\times P_a+P_b)$$

Introduction: Primes, a natural number greater than 1, except 1 and itself, can not be integer by other natural numbers except 1 and itself. In other words, the number has no other factors except 1 and itself; otherwise, it is called compound number. According to the basic theorem of arithmetic, every integer larger than 1 can be written either as a prime or as a product of a series of prime numbers; furthermore, if the order of these prime numbers in the product is not taken into account, the written form is unique. The smallest prime number is 2.

Abstract: prove logical thinking. Use the hypothesis theory to judge the unknown. In an infinite even number, there are only unknowns, a, b, c, d. Sole judgment; It is a prime number, or a compound number.

Arbitrary odd numbers: $【2N+1】 - P (P>2) = 2\times B(\text{Remainder } B, \text{ it's a prime, or it's a composite,})$ according to the judgment. If B equals a prime, there is no need to calculate it. Mathematical theory: B is an unknown. It can only be judged as a prime number, or a compound number. If it is a complex number, it can decompose the prime factor. So we can get prime numbers. Here, using computational logic theory: assuming that the remainder is a compound number, there is a transformation of the remainder, pushing the unknown to infinity
Let: any odd number $【(2N+1) \neq 2\times P_a+P_b】$ push the computational logic to infinity according to the analog computational logic: hypothesis theory. But any odd number $【2N+1】$ belongs to finite value, so it is contradictory. Mathematical theory: yes, or no. No: Negate all your questions.

Key words: hypothesis, prime number, compound number, decomposition prime factor. With reference to each other

According to the definition of prime number and complex number:

Lemma: Primes cannot be decomposed into prime factors.

Lemma: Composite number can decompose prime factors.

set up: (Natural Number N) $N > 1$

Arbitrary even number ($2N$)

Arbitrary odd number ($2N-1$)

$\therefore 2N+2=N(\text{Satisfying Integer Solution})$

\therefore Natural Number $N > 1$ (Prime number is odd)

set up: $2N+1 > 5$ Arbitrary odd number

$2N+1 \neq 2P_a+P_b$

set up: ($2N+1 \neq 2P_a+P_b$) Theoretical Conditions

Analog Synchronized Arithmetic Logic: $(2N+1)-2=R$

【R Be equal to: Prime number、 or、 Composite number】 In infinite numbers, we don't know whether R is a prime or a compound number. Here, we assume that complex number reasoning extends to infinite fields.

There is no contradiction between the arithmetic system and the hypothesis.

【Note: 2 does not participate in the following calculations if there is a prime number of 2 in alphabetical, composite decomposition】

$(2N+1)-2=R$ **【R Be equal to: Prime number、 or、 Composite number】**

If the remaining letters are prime, then **【 $2N+1=2 \times P_a+P_b$ 】**

hypothesis: R=Composite number: Prime factor decomposition $R=A \times B \times C$

Simulate synchronous arithmetic logic.

Extraction prime factor, A, B, C,

$$\left\{ \begin{array}{l} (2N+1)-A=2 \times E \text{ 【hypothesis: } E=\text{Composite number: Prime factor decomposition} \\ E=A^{n_1} \text{】} \\ (2N+1)-B=2 \times F \text{ 【hypothesis: } F=\text{Composite number: Prime factor decomposition} \\ F=B^{n_1} \text{】} \\ (2N+1)-C=2 \times G \text{ 【hypothesis: } G=\text{Composite number: Prime factor decomposition} \\ G=C^{n_1} \text{】} \end{array} \right.$$

Extraction prime factor, A_1 B_1 C_1

$$\left\{ \begin{array}{l} (2N+1)-A_1=2 \times E_1 \text{ 【hypothesis: } E_1=A^{n_2} \text{】} \\ (2N+1)-B_1=2 \times F_1 \text{ 【hypothesis: } F_1=B^{n_2} \text{】} \\ (2N+1)-C_1=2 \times G_1 \text{ 【hypothesis: } G_1=C^{n_2} \text{】} \end{array} \right.$$

Extraction prime factor, A_2 B_2 C_2

$$\left\{ \begin{array}{l} (2N+1)-A_2=2 \times E_2 \text{ 【hypothesis: } E_2=A^{n_3} \text{】} \\ (2N+1)-B_2=2 \times F_2 \text{ 【hypothesis: } F_2=B^{n_3} \text{】} \\ (2N+1)-C_2=2 \times G_2 \text{ 【hypothesis: } G_2=C^{n_3} \text{】} \end{array} \right.$$

Extraction prime factor, A_3 B_3 C_3

$$\left\{ \begin{array}{l} (2N+1)-A_3=2 \times E_3 \text{ 【hypothesis: } E_3=A^{n_4} \text{】} \\ (2N+1)-B_3=2 \times F_3 \text{ 【hypothesis: } F_3=B^{n_4} \text{】} \end{array} \right.$$

$$(2N+1)-C_3=2 \times G_3 \quad \text{【hypothesis: } G_3=C_4^n \text{】}$$

Extraction prime factor, A_4 、 B_4 、 C_4

.....Analog Arithmetic Logic (TY1)、

set up: **【 $(2N+1) \neq 2 \times P_a + P_b$ 】** according to arithmetic logic: there are only two alternative answers.

(一) Or arithmetic logic loop. The remainder letters are all compound numbers (cyclic prime numbers)

(二) On the contrary, arithmetic logic: Infinite non-circulate hypothesis (thus, Increase infinite, different prime numbers)

(一)、set up: **【 $(2N+1) \neq 2 \times P_a + P_b$ 】**

Analog cyclic arithmetic logic:

$$\begin{aligned} \because (2N+1) - 2 &= a \times b \\ \begin{cases} (2N+1) - a &= 2 \times 2^n \times c \\ (2N+1) - b &= 2 \times 2^m \end{cases} \\ \because \begin{cases} (2N+1) - 2 &= a \times b \\ (2N+1) - c &= 2 \times a^d \end{cases} \end{aligned}$$

Express, $(2N+1)$ With this arithmetic logic, (letter: a, b) remainder equals compound number According to (TY1) simulation arithmetic logic.

\therefore **【 $(2N+1) \neq 2 \times P_a + P_b$ 】**

(二)、set up: **【 $(2N+1) \neq 2 \times P_a + P_b$ 】**。

arithmetic logic: Infinite non-circulate hypothesis

(thus, Increase infinite, different prime numbers)

$$\because (2N+1) < \infty$$

The hypothesis is contradictory. On the contrary, according to the above arithmetic logic:

【 $(2N+1) = 2 \times P_a + P_b$ 】

hypothesis: $(2N-1 \neq 2P_1 + P_2)$ Under theoretical conditions, choice can only be (一), (TY1) arithmetic logic cycle

Then, simulate arithmetic.

$$(2N+1) - 2 = R \text{ (R=prime number、 or、 R=compound number) } \quad \text{【】}$$

hypothesis: R=compound number: Prime factor decomposition = $A \times B \times C$

If the remaining letters are prime, then **【 $2N+1 = 2 \times P_a + P_b$ 】**

Extraction prime factor, A、B、C、

$$\left\{ \begin{aligned} &(2N+1) - 2A = H \quad \text{【set up: H=compound number: Prime factor decomposition } H=L_1^n \text{】} \\ \textcircled{2} &(2N+1) - 2B = J \quad \text{【hypothesis: } J=M_1^n \text{】} \\ &(2N+1) - 2C = K \quad \text{【hypothesis: } K=W_1^n \text{】} \end{aligned} \right.$$

Extraction Prime Factors, L_1 、 M_1 、 W_1

$$\textcircled{1} \begin{cases} (2N+1) - L_1 = 2H_1 \div 2 \quad \text{【hypothesis: } H_1=L_2^n \text{】} \\ (2N+1) - M_1 = 2J_1 \div 2 \quad \text{【hypothesis: } J_1=M_2^n \text{】} \\ (2N+1) - W_1 = 2K_1 \div 2 \quad \text{【hypothesis: } K_1=W_2^n \text{】} \end{cases}$$

Extraction prime factor, L_2 、 M_2 、 W_2

$$\textcircled{2} \begin{cases} (2N+1) - 2L_2 = H_2 \quad \text{【hypothesis: } H_2=L_3^n \text{】} \\ (2N+1) - 2M_2 = J_2 \quad \text{【hypothesis: } J_2=M_3^n \text{】} \\ (2N+1) - 2W_2 = K_2 \quad \text{【hypothesis: } K_2=W_3^n \text{】} \end{cases}$$

Extraction prime factor、 L_3 、 M_3 、 W_3

$$\textcircled{1} \begin{cases} (2N+1) - L_3 = 2H_3 \div 2 \text{ 【hypothesis: } K_3 = L^n_4 \text{】} \\ (2N+1) - M_3 = 2J_3 \div 2 \text{ 【hypothesis: } K_3 = M^n_4 \text{】} \\ (2N+1) - W_3 = 2K_3 \div 2 \text{ 【hypothesis: } K_3 = W^n_4 \text{】} \end{cases}$$

Extraction prime factor、 L_4 、 M_4 、 W_4

$$\textcircled{2} \begin{cases} (2N+1) - 2L_4 = H_4 \text{ 【hypothesis: } H_4 = L^n_5 \text{】} \\ (2N+1) - 2M_4 = J_4 \text{ 【hypothesis: } J_4 = M^n_5 \text{】} \\ (2N+1) - 2W_4 = K_4 \text{ 【hypothesis: } K_4 = W^n_5 \text{】} \end{cases}$$

Extraction prime factor、 L_5 、 M_5 、 W_5

$$\textcircled{1} \begin{cases} (2N+1) - L_5 = 2H_5 \div 2 \text{ 【hypothesis: } H_5 = L^n_6 \text{】} \\ (2N+1) - M_5 = 2J_5 \div 2 \text{ 【hypothesis: } J_5 = M^n_6 \text{】} \\ (2N+1) - W_5 = 2K_5 \div 2 \text{ 【hypothesis: } K_5 = W^n_6 \text{】} \end{cases}$$

Extraction prime factor、 L_6 、 M_6 、 W_6

..... $\textcircled{2}\textcircled{1}\textcircled{2}\textcircled{1}$ Analog Arithmetic Logic (TY2)。

set up: 【 $(2N+1) \neq 2 \times P_a + P_b$ 】 according to arithmetic logic: there are only two alternative answers.

(三)、Either, loop arithmetic logic. The remainder (unknowns, letters) are compound Numbers

(四)、On the contrary, arithmetic logic: Infinite non-circulate hypothesis

(thus, Increase infinite, different prime numbers)

(三)、set up: : Arithmetic Logic Cycle (Cyclic Prime)

Expressions, $(2N+1)$ According to arithmetic logic, (letter) remainder equals compound number

$(2N+1 \neq 2P_a + P_b)$

(四)、On the contrary, arithmetic logic, infinitely increasing assumptions

(thus: infinitely increasing, different prime numbers)

Expressing, $2N+1$ uses this arithmetic logic to determine that the remainder is equal to the compound number. $(2N+1 \neq 2P_a + P_b)$ 。

(So: infinite increase, different prime numbers)

\therefore Arbitrarily $(2N+1) < \infty$

\therefore Assuming contradictions, according to the above arithmetic logic 【 $(2N+1) = 2 \times P_a + P_b$ 】

Assumptions: $2N + 1 \neq 2P_a + P_b$ (under theoretical conditions), selection (一), cyclic arithmetic (cyclic prime) and (三) cyclic arithmetic logic

The two arithmetic logic and circular hypothesis problems are merged into one problem for mutual reference judgment and reasoning.

(TY2) and (TY1) prime cyclic arithmetic, (TY2) refer to (TY1) cyclic arithmetic logic. Each step of judgment

Analog (TY1) Cyclic Arithmetic Logic

$\therefore (2N+1) - 2 = a \times b$

$$\begin{cases} (2N+1) - a = 2 \times 2^n \times c \\ (2N+1) - b = 2 \times 2^m \end{cases}$$

$$\begin{cases} (2N+1) - 2 = a^a \times b^b \\ (2N+1) - c = 2 \times a^d \end{cases}$$

\therefore According to analog arithmetic logic 【 $(2N+1) \neq 2 \times P_a + P_b$ 】

(TY2) Reference (TY1) Arithmetic Logic

$(2N+1) - 2 = a \times b$

(TY2) Arithmetic Logic Hypothesis:

$$\begin{cases} (2N+1) - 2 \times a = 2 \times 2^n \times c \\ (2N+1) - 2 \times b = 2 \times 2^n \end{cases}$$

The equation holds.

$$\begin{cases} (2N+1) - 2 \times a \text{ 【Odd number—Even numbers= Odd number】} \\ (2N+1) - 2 \times b \text{ 【Odd number—Even numbers= Odd number】} \end{cases}$$

Hypothetical contradiction.

$$\begin{cases} (2N+1) - 2 \times 2 = a^a \times b^b \\ (2N+1) - 2 \times c = a^d \\ (2N+1) - 2 = a^a \times b^b \\ (2N+1) - 2 \times 2 = a^c \times b^d \end{cases}$$

Solution: $2 = a^{c-a} \times b^{b-d}$

The smallest prime is equal to 2

$$\therefore 2 < a^{c-a} \times b^{b-d}$$

Hypothetical contradiction.

$$\begin{cases} (2N+1) - c = 2 \times a^d \\ (2N+1) - 2 \times c = a^e \\ c = a^d(a^{e-d} - 2) \end{cases}$$

$\therefore c =$ (Primes cannot be decomposed)

$\therefore a =$ Scores (not in the set of integers)

Assumptions are contradictory. On the contrary, according to (TY2) arithmetic logic, assumptions are infinitely increased (thus: infinitely increased, different primes)

$$\therefore \text{Arbitrarily: } (2N+1) < \infty$$

\therefore Assuming contradiction, contrary, according to the above arithmetic logic 【 $(2N+1) = 2 \times P_a + P_b$ 】

In the infinite field, we cannot judge one by one. If one judgement belongs to exhaustive law, it belongs to pseudoscience.

Simulated (TY1) cyclic arithmetic logic is performed with S-term 【Note: 2. No calculation in it.】

$$\begin{aligned} (2N+1) - 2 &= A \times B \\ \begin{cases} (2N+1) - A = 2 \times C^c & 1 \\ (2N+1) - B = 2 \times D^d \end{cases} \\ \begin{cases} (2N+1) - C = 2 \times C_1^{c_1} & 2 \\ (2N+1) - D = 2 \times D_1^{d_1} \end{cases} \\ \begin{cases} (2N+1) - C_1 = 2 \times C_2^{c_2} & 3 \\ (2N+1) - D_1 = 2 \times D_2^{d_2} \end{cases} \\ \begin{cases} (2N+1) - C_2 = 2 \times C_3^{c_3} & 4 \\ (2N+1) - D_2 = 2 \times D_3^{d_3} \end{cases} \\ \begin{cases} (2N+1) - C_3 = 2 \times C_4^{c_4} & 5 \\ (2N+1) - D_3 = 2 \times D_4^{d_4} \end{cases} \\ \begin{cases} (2N+1) - C_4 = 2 \times C_5^{c_5} & 6 \\ (2N+1) - D_4 = 2 \times D_5^{d_5} \end{cases} \\ \dots\dots\dots S \text{ term} \\ \begin{cases} (2N+1) - C_s = 2 \times A^a & 1 \\ (2N+1) - D_s = 2 \times B^b \end{cases} \\ \begin{cases} (2N+1) - A = 2 \times C^c & 2 \\ (2N+1) - B = 2 \times D^d \end{cases} \end{aligned}$$

Simulate the arithmetic logic according to the above (the remainder is compound)

$$\therefore \text{【 } (2N+1) = 2 \times P_a + P_b \text{】}$$

Then, (TY2) arithmetic logic mutual reference (TY1) circular arithmetic logic. All prime numbers.

(TY1) Primes per step, simulating arithmetic logic according to (TY2)

Here is item S.

$$\begin{array}{l}
 \left\{ \begin{array}{l} (2N+1) - 2C_1 = \mathbf{[H_1]} \\ (2N+1) - 2D_1 = \mathbf{[M_1]} \end{array} \right. \quad (\text{WY2}) \textcircled{2}\textcircled{1}\textcircled{2}\textcircled{1}\dots\dots \\
 \left\{ \begin{array}{l} (2N+1) - 2C_2 = \mathbf{[H_2]} \\ (2N+1) - 2D_2 = \mathbf{[M_3]} \end{array} \right. \quad (\text{WY2}) \textcircled{2}\textcircled{1}\textcircled{2}\textcircled{1}\dots\dots \\
 \left\{ \begin{array}{l} (2N+1) - 2C_3 = \mathbf{[H_3]} \\ (2N+1) - 2D_3 = \mathbf{[M_3]} \end{array} \right. \quad (\text{WY2}) \textcircled{2}\textcircled{1}\textcircled{2}\textcircled{1}\dots\dots \\
 \left\{ \begin{array}{l} (2N+1) - 2C_4 = \mathbf{[H_4]} \\ (2N+1) - 2D_4 = \mathbf{[M_4]} \end{array} \right. \quad (\text{WY2}) \textcircled{2}\textcircled{1}\textcircled{2}\textcircled{1}\dots\dots \\
 \left\{ \begin{array}{l} (2N+1) - 2C_5 = \mathbf{[H_5]} \\ (2N+1) - 2D_5 = \mathbf{[M_5]} \end{array} \right. \quad (\text{WY2}) \textcircled{2}\textcircled{1}\textcircled{2}\textcircled{1}\dots\dots \\
 \dots\dots\dots (\text{WY2}) \textcircled{2}\textcircled{1}\textcircled{2}\textcircled{1}\dots\dots \\
 \left\{ \begin{array}{l} (2N+1) - 2C_s = \mathbf{[H_1]} \\ (2N+1) - 2D_s = \mathbf{[M_1]} \end{array} \right. \quad (\text{WY2}) \textcircled{2}\textcircled{1}\textcircled{2}\textcircled{1}\dots\dots
 \end{array}$$

[H₁] ~~~ [H_s] [M₁] ~~~ [M_s] They are all complex numbers., Prime factor decomposition hypothesis: **[H₁] ~~~ [H_s] [M₁] ~~~ [M_s]** (TY1) The same prime factor.

that:

$$\begin{array}{l}
 \left\{ \begin{array}{l} (2N+1) \div 2 - C_1 = \mathbf{[H_1]} \\ (2N+1) \div 2 - D_1 = \mathbf{[M_1]} \end{array} \right. \\
 \left\{ \begin{array}{l} (2N+1) \div 2 - C_2 = \mathbf{[H_2]} \\ (2N+1) \div 2 - D_2 = \mathbf{[M_2]} \end{array} \right. \\
 \left\{ \begin{array}{l} (2N+1) \div 2 - C_3 = \mathbf{[H_3]} \\ (2N+1) \div 2 - D_3 = \mathbf{[M_3]} \end{array} \right. \\
 \left\{ \begin{array}{l} (2N+1) \div 2 - C_4 = \mathbf{[H_4]} \\ (2N+1) \div 2 - D_4 = \mathbf{[M_4]} \end{array} \right. \\
 \left\{ \begin{array}{l} (2N+1) \div 2 - C_5 = \mathbf{[H_5]} \\ (2N+1) \div 2 - D_5 = \mathbf{[M_5]} \end{array} \right. \\
 \dots\dots\dots \\
 \left\{ \begin{array}{l} (2N+1) \div 2 - C_s = \mathbf{[H_s]} \\ (2N+1) \div 2 - D_s = \mathbf{[M_s]} \end{array} \right.
 \end{array}$$

The same arithmetic, the same logic.

(2N+1) ÷2 (It's a fraction, not an integer.)

Assuming contradictions,

Then, (TY1) Arithmetic Logic Cycle, (TY2) Arithmetic Logic Increases Infinitely
 $(2N+1) < \infty$

Assuming contradictions, according to arithmetic logic (TY2) **[(2N+1) = 2×P_a + P_b]**

Reference document:

Euclid. Infinite theorems of prime numbers.Cited edition of scientific research

Abstract Hypothesis: Simulating Basic Arithmetic Logic, Judging Reasoning and Hypothesis Contradictions

Condition <integer theory>

Abstract hypothesis:

Let: the number of primes be finite.

The order from small to large is $P_1, P_2, P_3, \dots, P_n$

Simulated Basic Arithmetic Logic: Multiplication from Small to Large

$$P_1 \times P_2 \times P_3 \times \dots \times P_n = N$$

$$2 \times 3 \times 5 \times 7 \times \dots \times P_n = N$$

that: $(N+1)$

$$N+1 > P_n$$

Judgement reasoning:

If $N+1$ is a composite number,

set up: $N+1=W$ **[X]**

set up: $W=P_1, P_2, P_3, \dots, P_n$ (Arbitrary prime number)

$$(N+1) \div W$$

$N \div W$ (Satisfying Integer Solution)

$1 \div W$ (Fraction)

hypothesis: $(N+1) \div W=X$

If **[$(N+1) \div W=X$]** hypothesis holds

$\therefore X$ is a fraction, not an integer set.

\therefore Conversely, the hypothesis is not valid.

$(N+1)$ complex or prime

$(N+1)$ Factor obtained by decomposition, It's certainly not in the assumptions of $P_1, P_2, P_3, \dots, P_n$

There are other prime numbers in addition to the finite prime numbers assumed.

So the original assumption is not valid. That is to say, there are infinite primes.

Making Web Sites Based on Thesis Simulated Computing Logic:

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odd number = 2^* (prime number) + prime number

$$11111111111111111111 = 2 * 444444444444379 + 11022222222222353$$

$$11111111111111111113 = 2 * 11 + 11111111111111091$$

$$11111111111111111115 = 2 * 333667 + 111111111110443781$$

$$11111111111111111117 = 2 * 13 + 11111111111111091$$

$$11111111111111111119 = 2 * 5555555555519841 + 71437$$

$$11111111111111111121 = 2 * 555555555547947 + 15227$$

$$11111111111111111123 = 2 * 5555555548482753 + 14145617$$

$$11111111111111111125 = 2 * 5555554914907011 + 1281297103$$

$$11111111111111111127 = 2 * 55535945405341813 + 39220300427501$$

$$11111111111111111129 = 2 * 43 + 11111111111111043$$

$$\begin{aligned}
9 &= 2 \times 2 + 5 \\
99 &= 2 \times 5 + 89 \\
999 &= 2 \times 199 + 601 \\
9999 &= 2 \times 13 + 9973 \\
99999 &= 2 \times 5 + 99989 \\
999999 &= 2 \times 49999 + 900001 \\
9999999 &= 2 \times 4999453 + 1093 \\
99999999 &= 2 \times 5 + 99999989 \\
999999999 &= 2 \times 499679693 + 640613 \\
9999999999 &= 2 \times 4409 + 9999991181 \\
99999999999 &= 2 \times 1483 + 99999997033 \\
999999999999 &= 2 \times 5 + 99999999989 \\
9999999999999 &= 2 \times 499999999769 + 461 \\
99999999999999 &= 2 \times 24999999384823 + 50000001230353 \\
999999999999999 &= 2 \times 5 + 9999999999989 \\
9999999999999999 &= 2 \times 41 + 99999999999917 \\
99999999999999999 &= 2 \times 216273151 + 9999999567453697 \\
999999999999999999 &= 2 \times 5 + 999999999999989 \\
33 &= 2 \times 2 + 29 \\
333 &= 2 \times 47 + 239 \\
3333 &= 2 \times 2 + 3329 \\
33333 &= 2 \times 2 + 33329 \\
333333 &= 2 \times 166643 + 47 \\
3333333 &= 2 \times 166657 + 19 \\
33333333 &= 2 \times 2 + 3333329 \\
333333333 &= 2 \times 3333331 + 26666671 \\
3333333333 &= 2 \times 166666637 + 59 \\
33333333333 &= 2 \times 2 + 333333329 \\
333333333333 &= 2 \times 17 + 3333333299 \\
3333333333333 &= 2 \times 41 + 33333333251 \\
33333333333333 &= 2 \times 7 + 33333333319 \\
333333333333333 &= 2 \times 165938864628821 + 1455604075691 \\
3333333333333333 &= 2 \times 103 + 3333333333127 \\
33333333333333333 &= 2 \times 16666666666661 + 11
\end{aligned}$$

抽象假设，模拟同数算术逻辑，判断推理和假设矛盾

大于 5 的奇数等于 2 加一个奇素数
或者大于 5 的奇数等于 2 个相同素数加一个奇素数
(简称: $2N+1=2+P_a$ 、或、 $2N+1=2 \times P_a+P_b$)

反证法 (整数理论, 哲学)

素数。一个大于 1 的自然数，除了 1 和它自身外，不能被其他自然数整除的数叫做质数；否则称为合数。

摘要：证明逻辑思路，假设任意奇数：【 $(2N+1) \neq 2 \times P_a + P_b$ 】，那么根据模拟计算逻辑。

用假设判断未知。在无穷大数里面，未知数、a、b、c、d。只能依据判断；它素数、或合数。

当：【 $2N+1$ 】 $-P$ ($P>2$) $=2 \times B$ (余项 B、它是素数、或者、它是合数，) 根判断据。如果 B 等于质数，就不需要计算。但是 B 是一个未知数，只能根据判断它是素数、或合数。如果是合数，它就可以分解质因数。这样我们就可以得到素数。这里，用计算逻辑理论：假设余项都是合数，用余项存在变化把未知数推到无穷大。求：【 $(2N+1) = 2 \times P_a + P_b$ 】

根据计算逻辑，假设把任意【 $(2N+1) \neq 2 \times P_a + P_b$ 】计算推进到无穷，而任意【 $2N+1$ 】属于有限，因此得到矛盾。

反证法：要么是、要么否、否：否定一切你提出的问题

关键词：假设、素数、合数、分解质因数。相互参照

数学理论判断：这是什么数学理论。

问：那么这是一遍整数理论、还是分数理论。

假设它是一遍分数理论，相反用分数判断一个整数，我们不能判断它是素数、或者复合数
 \therefore 它是一遍整数理论

根据素数和合数定义

定理：素数不可以分解质因数。

定理：复合数可以分解质因数。

根据素数、合数、定义。

设：(自然数 N) $N>1$

任意偶数 $=2N$

任意奇数 $=2N-1$

$\therefore 2N+2=N$ (满足整数解)

\therefore 自然数 $N>1$ (偶数里面没有素数)

设： $2N+1>5$ 的任意奇数

假设：【 $2N+1 \neq 2 \times P_a + P_b$ 】情况下

模拟同步算术逻辑： $(2N+1) - 2 = R$

【R：要么是素数，要么是合数】在无限大的数字里面，我们不懂得 R 是素数或者是合数，这里假设合数推理，延伸到无限大领域。

假设与算术系统并不矛盾 (欧几里得反证法科学研究改写版、可以解答)

【注：如果字母，合数分解，存在 2 的素数，那么 2 不参与下列计算】

模拟同步算术逻辑。

$(2N+1) - 2 = R$ (R 素数、或、复合数) 【无限大数学里面，R 属于未知数我们不知道素数或者是合数】

如果余项字母是质数、那么 【 $2N+1 = 2 \times P_a + P_b$ 】

重要余项的数在变化

假设：R=复合数，合数分解 $R=A \times B \times C$

$$\left\{ \begin{array}{l} (2N+1) - A = 2 \times E \text{ 【假设：E=合数：E 因式分解 } E = A^{n_1} \text{】} \\ (2N+1) - B = 2 \times F \text{ 【假设：F=合数：F 因式分解 } F = B^{n_1} \text{】} \end{array} \right.$$

$(2N+1)-C=2\times G$ 【假设: $G=$ 合数: G 因式分解 $G=C^n_1$ 】

抽取素因数 $A_1 B_1 C_1$

$$\begin{cases} (2N+1)-A_1=2\times E_1 \text{ 【假设: } E_1=A^n_2 \text{】} \\ (2N+1)-B_1=2\times F_1 \text{ 【假设: } F_1=B^n_2 \text{】} \\ (2N+1)-C_1=2\times G_1 \text{ 【假设: } G_1=C^n_2 \text{】} \end{cases}$$

抽取素因数 $A_2 B_2 C_2$

$$\begin{cases} (2N+1)-A_2=2\times E_2 \text{ 【假设: } E_2=A^n_3 \text{】} \\ (2N+1)-B_2=2\times F_2 \text{ 【假设: } F_2=B^n_3 \text{】} \\ (2N+1)-C_2=2\times G_2 \text{ 【假设: } G_2=C^n_3 \text{】} \end{cases}$$

抽取素因数 $A_3 B_3 C_3$

$$\begin{cases} (2N+1)-A_3=2\times E_3 \text{ 【假设: } E_3=A^n_4 \text{】} \\ (2N+1)-B_3=2\times F_3 \text{ 【假设: } F_3=B^n_4 \text{】} \\ (2N+1)-C_3=2\times G_3 \text{ 【假设: } G_3=C^n_4 \text{】} \end{cases}$$

抽取素因数 A_4 、 B_4 、 C_4

..... 模拟算术逻辑 (TY1)、

算术逻辑, 设: 【 $(2N+1) \neq 2\times P_a+P_b$ 】条件下。只有两个选择答案。

(一) 要么、算术逻辑循环。余项字母都是合数 (算术循环逻辑)

(二) 相反、算术逻辑不循环, 无限增加假设 (于是: 增加无穷个, 不相同素数)

(一)、假设: 【 $(2N+1) \neq 2\times P_a+P_b$ 】

模拟循环算术逻辑:

$$\begin{aligned} \because (2N+1)-2 &= a \times b \\ \begin{cases} (2N+1)-a &= 2 \times 2^n \times c \\ (2N+1)-b &= 2 \times 2^m \end{cases} \\ \because \begin{cases} (2N+1)-2 &= a \times b \\ (2N+1)-c &= 2 \times a^d \end{cases} \end{aligned}$$

\therefore 根据上面算术模拟逻辑, 【 $(2N+1) \neq 2\times P_a+P_b$ 】

表达 $(2N+1)$ 用这个算术逻辑, (字母: a 、 b) 余项判断等于合数

(三)、设: 【 $(2N+1) \neq 2\times P_a+P_b$ 】。

算术逻辑, 无限增加假设 (于是: 无限增加, 不同素数)

$$\because (2N+1) < \infty$$

假设矛盾。相反, 根据上面算术逻辑 【 $(2N+1) = 2\times P_a+P_b$ 】

假设: $(2N-1 \neq 2P_1+P_2)$ 条件下、选择只能 (一), (TY1) 算术逻辑循环

再进行, 模拟算术。

$(2N+1)-2=R$ ($R=$ 素数、或、 $R=$ 复合数) 【】

假设: $R=$ 复合数: 合数分解= $A \times B \times C$

如果余项字母是质数、那么 【 $2N+1=2\times P_a+P_b$ 】

抽取素因数、 A 、 B 、 C 、

$$\textcircled{2} \begin{cases} (2N+1)-2A=H \text{ 【假设: } H= \text{合数: 合数分解 } H=L^n_1 \text{】} \\ (2N+1)-2B=J \text{ 【假设: } J= \text{合数: 合数分解 } J=M^n_1 \text{】} \end{cases}$$

$$(2N+1) - 2C = K \text{ 【假设: } K = \text{合数: 合数分解 } K = W^n_1 \text{】}$$

抽取素因数、 L_1 、 M_1 、 W_1

$$\textcircled{1} \begin{cases} (2N+1) - L_1 = 2H_1 \div 2 \text{ 【假设: } H_1 = L^n_2 \text{】} \\ (2N+1) - M_1 = 2J_1 \div 2 \text{ 【假设: } J_1 = M^n_2 \text{】} \\ (2N+1) - W_1 = 2K_1 \div 2 \text{ 【假设: } K_1 = W^n_2 \text{】} \end{cases}$$

抽取素因数、 L_2 、 M_2 、 W_2

$$\textcircled{2} \begin{cases} (2N+1) - 2L_2 = H_2 \text{ 【假设: } H_2 = L^n_3 \text{】} \\ (2N+1) - 2M_2 = J_2 \text{ 【假设: } J_2 = M^n_3 \text{】} \\ (2N+1) - 2W_2 = K_2 \text{ 【假设: } K_2 = W^n_3 \text{】} \end{cases}$$

抽取素因数、 L_3 、 M_3 、 W_3

$$\textcircled{1} \begin{cases} (2N+1) - L_3 = 2H_3 \div 2 \text{ 【假设: } H_3 = L^n_4 \text{】} \\ (2N+1) - M_3 = 2J_3 \div 2 \text{ 【假设: } J_3 = M^n_4 \text{】} \\ (2N+1) - W_3 = 2K_3 \div 2 \text{ 【假设: } K_3 = W^n_4 \text{】} \end{cases}$$

抽取素因数、 L_4 、 M_4 、 W_4

$$\textcircled{2} \begin{cases} (2N+1) - 2L_4 = H_4 \text{ 【假设: } H_4 = L^n_5 \text{】} \\ (2N+1) - 2M_4 = J_4 \text{ 【假设: } J_4 = M^n_5 \text{】} \\ (2N+1) - 2W_4 = K_4 \text{ 【假设: } K_4 = W^n_5 \text{】} \end{cases}$$

抽取素因数、 L_5 、 M_5 、 W_5

$$\textcircled{1} \begin{cases} (2N+1) - L_5 = 2H_5 \div 2 \text{ 【假设: } H_5 = L^n_6 \text{】} \\ (2N+1) - M_5 = 2J_5 \div 2 \text{ 【假设: } J_5 = M^n_6 \text{】} \\ (2N+1) - W_5 = 2K_5 \div 2 \text{ 【假设: } K_5 = W^n_6 \text{】} \end{cases}$$

抽取素因数、 L_6 、 M_6 、 W_6

..... $\textcircled{2} \textcircled{1} \textcircled{2} \textcircled{1}$ 模拟算术 (TY2)

算术逻辑, 设: $[(2N+1) \neq 2 \times P_a + P_b]$ 条件下。只能两个选择

(三)、要么、循环算术逻辑。余项字母都是合数 (素数循环)

(四)、相反, 算术逻辑: 无限非循环假设 (因此, 增加了无穷多的不同素数)

(三)、假设: 算术逻辑循环 (素数循环)

表达、 $(2N+1)$ 用这个算术逻辑, (字母) 余项判断等于复合数

$$\therefore (2N+1) \neq 2P_a + P_b$$

(四)、相反, 算术逻辑: 无限非循环假设 (因此, 增加了无穷多的不同素数)

表达, $2N+1$ 用这个算术逻辑, (字母) 余项判断等于复合数 $(2N+1) \neq 2P_a + P_b$ 。

(因此, 增加了无穷多的不同素数)

$$\therefore \text{任意 } (2N+1) < \infty$$

$$\therefore \text{假设矛盾, 相反根据上面算术逻辑 } [(2N+1) = 2 \times P_a + P_b]$$

假设: $2N+1 \neq 2P_a + P_b$ (理论条件下)、选择 (一)、循环算术 (素数循环) 和 (三) 循环算术逻辑

将两个算术逻辑, 循环假设问题, 合并成一个问题, 进行相互参照判断推理。

(TY2) 和 (TY1) 俩式素数循环算术, (TY2) 参照 (TY1) 算术逻辑循环。每一步判段

模拟 (TY1) 循环算术逻辑

$$\therefore (2N+1) - 2 = a \times b$$

$$\begin{cases} (2N+1) - a = 2 \times 2^n \times c \\ (2N+1) - b = 2 \times 2^m \end{cases}$$

$$\begin{cases} (2N+1) - 2 = a^a \times b^b \\ (2N+1) - c = 2 \times a^d \end{cases}$$

∴根据算术逻辑【(2N+1) ≠ 2×P_a+P_b】

(TY2) 参照 (TY1) 算术逻辑

$$(2N+1) - 2 = a \times b$$

(TY2) 算术逻辑假设:

$$\begin{cases} (2N+1) - 2 \times a = 2 \times 2^n \times c \\ (2N+1) - 2 \times b = 2 \times 2^n \end{cases}$$

等式成立。

$$\begin{cases} (2N+1) - 2 \times a \text{ 【奇数 - 偶数 = 奇数】} \\ (2N+1) - 2 \times b \text{ 【奇数 - 偶数 = 奇数】} \end{cases}$$

假设矛盾。

$$\begin{cases} (2N+1) - 2 \times 2 = a^a \times b^b \\ (2N+1) - 2 \times c = a^d \end{cases}$$

$$\begin{cases} (2N+1) - 2 = a^a \times b^b \\ (2N+1) - 2 \times 2 = a^c \times b^d \end{cases}$$

解得: $2 = a^{c-a} \times b^{b-d}$

最小的素数等于 2

$$\therefore 2 < a^{c-a} \times b^{b-d}$$

假设矛盾

$$\begin{cases} (2N+1) - c = 2 \times a^d \\ (2N+1) - 2 \times c = a^e \end{cases}$$

解得: $c = a^d(a^{e-d} - 2)$

∴c=素数不能分解

∴a=分数 (不在整数集合里)

假设矛盾, 相反根据 (TY2) 算术逻辑, 增加新的素因数 (于是: 无限增加, 不同素数)

∴任意 (2N+1) < ∞

∴假设矛盾, 相反根据上面算术逻辑【(2N+1) = 2×P_a+P_b】

模拟 (TY1) 循环算术逻辑用 S 项进行【注: 2 在里面就不参加计算】.....

$$(2N+1) - 2 = A \times B$$

$$\begin{cases} (2N+1) - A = 2 \times C^c & 1 \\ (2N+1) - B = 2 \times D^d \end{cases}$$

$$\begin{cases} (2N+1) - C = 2 \times C^{c_1} & 2 \\ (2N+1) - D = 2 \times D^{d_1} \end{cases}$$

$$\begin{cases} (2N+1) - C_1 = 2 \times C^{c_2} & 3 \\ (2N+1) - D_1 = 2 \times D^{d_2} \end{cases}$$

$$\begin{cases} (2N+1) - C_2 = 2 \times C^{c_3} & 4 \\ (2N+1) - D_2 = 2 \times D^{d_3} \end{cases}$$

$$\begin{cases} (2N+1) - C_3 = 2 \times C^{c_4} & 5 \\ (2N+1) - D_3 = 2 \times D^{d_4} \end{cases}$$

$$\begin{cases} (2N+1) - C_4 = 2 \times C^{c_5} & 6 \\ (2N+1) - D_4 = 2 \times D^{d_5} \end{cases}$$

..... S 项

$$\begin{cases} (2N+1) - C_s = 2 \times A^a & 1 \\ (2N+1) - D_s = 2 \times B^b \end{cases}$$

$$\begin{cases} (2N+1) - A = 2 \times C^c & 2 \\ (2N+1) - B = 2 \times D^d \end{cases}$$

根据上面模拟算术逻辑 (余项都是合数)

$$\therefore \text{【 } (2N+1) = 2 \times P_a + P_b \text{】}$$

那么、(TY2) 算术逻辑【相互参照】(TY1) 算术逻辑循环。全部素数。

(TY1) 每步素数, 根据 (TY2) 算术逻辑

那么把、 $(2N+1) - 2P_a = W^w$

【W 属于 (TY1) 素数】

$(2N+1) - 2W$ 【根据 (TY2) 算术逻辑再进行, 模拟逻辑计算】

那么这里有: S 项

$$\begin{cases} (2N+1) - 2C_1 = \text{【H}_1\text{】} & (\text{WY2}) \text{ ②①②①} \dots\dots \\ (2N+1) - 2D_1 = \text{【M}_1\text{】} & \end{cases}$$

$$\begin{cases} (2N+1) - 2C_2 = \text{【H}_2\text{】} & (\text{WY2}) \text{ ②①②①} \dots\dots \\ (2N+1) - 2D_2 = \text{【M}_3\text{】} & \end{cases}$$

$$\begin{cases} (2N+1) - 2C_3 = \text{【H}_3\text{】} & (\text{WY2}) \text{ ②①②①} \dots\dots \\ (2N+1) - 2D_3 = \text{【M}_3\text{】} & \end{cases}$$

$$\begin{cases} (2N+1) - 2C_4 = \text{【H}_4\text{】} & (\text{WY2}) \text{ ②①②①} \dots\dots \\ (2N+1) - 2D_4 = \text{【M}_4\text{】} & \end{cases}$$

..... (WY2) ②①②①.....

$$\begin{cases} (2N+1) - 2C_5 = \text{【H}_5\text{】} & (\text{WY2}) \text{ ②①②①} \dots\dots \\ (2N+1) - 2D_5 = \text{【M}_5\text{】} & \end{cases}$$

$$\begin{cases} (2N+1) - 2C_s = \text{【H}_1\text{】} & (\text{WY2}) \text{ ②①②①} \dots\dots \\ (2N+1) - 2D_s = \text{【M}_1\text{】} & \end{cases}$$

【H₁】 ~~~ 【H_s】 【M₁】 ~~~ 【M_s】 全部都是合数, 分解素因数。

假设 (TY1) 质因数和 【H₁】 ~~~ 【H_s】 【M₁】 ~~~ 【M_s】 全部相同。

那么有,

$$\begin{cases} (2N+1) \div 2 - C_1 = \text{【H}_1\text{】} \\ (2N+1) \div 2 - D_1 = \text{【M}_1\text{】} \end{cases}$$

$$\begin{cases} (2N+1) \div 2 - C_2 = \text{【H}_2\text{】} \\ (2N+1) \div 2 - D_2 = \text{【M}_2\text{】} \end{cases}$$

$$\begin{cases} (2N+1) \div 2 - C_3 = \text{【H}_3\text{】} \\ (2N+1) \div 2 - D_3 = \text{【M}_3\text{】} \end{cases}$$

$$\begin{cases} (2N+1) \div 2 - C_4 = \text{【H}_4\text{】} \\ (2N+1) \div 2 - D_4 = \text{【M}_4\text{】} \end{cases}$$

..... (WY2) ②①②①.....

$$\begin{cases} (2N+1) \div 2 - C_5 = \text{【H}_s\text{】} \\ (2N+1) \div 2 - D_5 = \text{【M}_s\text{】} \end{cases}$$

$$\begin{cases} (2N+1) \div 2 - C_s = \text{【H}_1\text{】} \\ (2N+1) \div 2 - D_s = \text{【M}_1\text{】} \end{cases}$$

$(2N+1) \div 2$ (不满足整数解)

假设矛盾,

相反、根据算术逻辑逻辑 (TY2) 增加不相同素数 (于是、无限增加, 不相同素数)

\therefore 任意 $(2N+1) <$ 无限大

\therefore 假设矛盾, 信反根据算术基本逻辑 (TY2) 【 $(2N+1) = 2 \times P_a + P_b$ 】